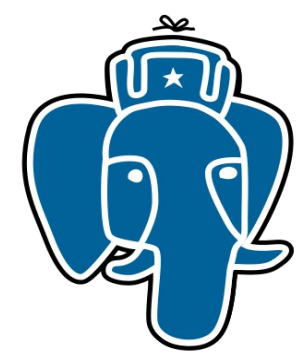


Schema-less PostgreSQL

Current and Future

September, 2014, Tokyo, Japan

Oleg Bartunov, SAI MSU

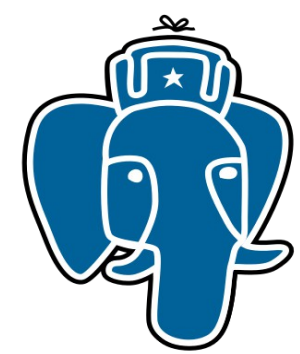


Oleg Bartunov, Teodor Sigaev

- Locale support
- Extendability (indexing)
 - GiST (KNN), GIN, SP-GiST
- Full Text Search (FTS)
- Jsonb, VODKA
- Extensions:
 - intarray
 - pg_trgm
 - ltree
 - hstore
 - plantuner



<https://www.facebook.com/oleg.bartunov>
obartunov@gmail.com, teodor@sigaev.ru
<https://www.facebook.com/groups/postgresql/>

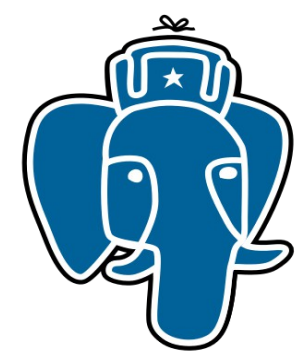


Alexander Korotkov

- Indexed regexp search
- GIN compression & fast scan
- Fast GiST build
- Range types indexing
- Split for GiST

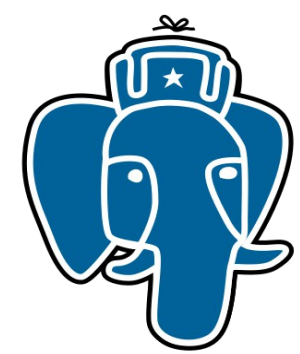


akeorotkov@gmail.com



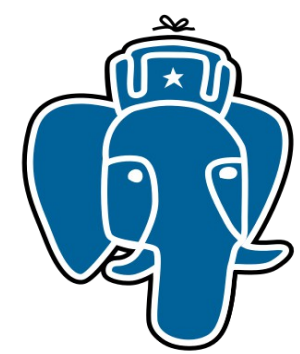
Agenda

- The problem
- Hstore
- Introduction to jsonb indexing
- JQuery - Jsonb Query Language
- Exercises on jsonb GIN opclasses with JQuery support
- VODKA access method



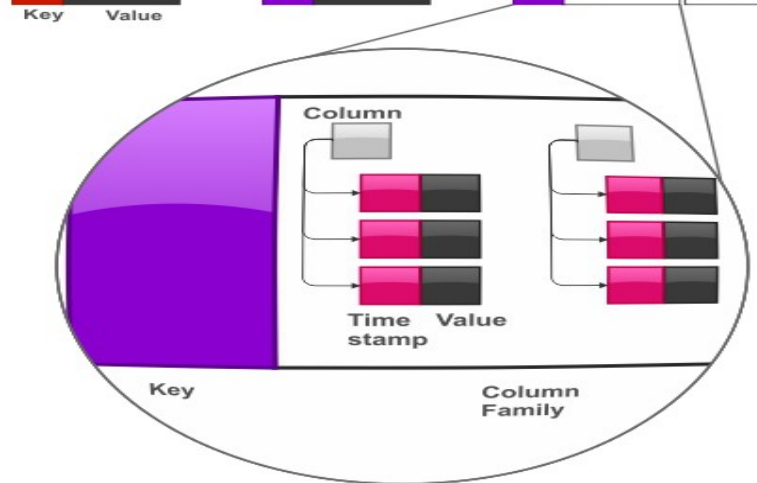
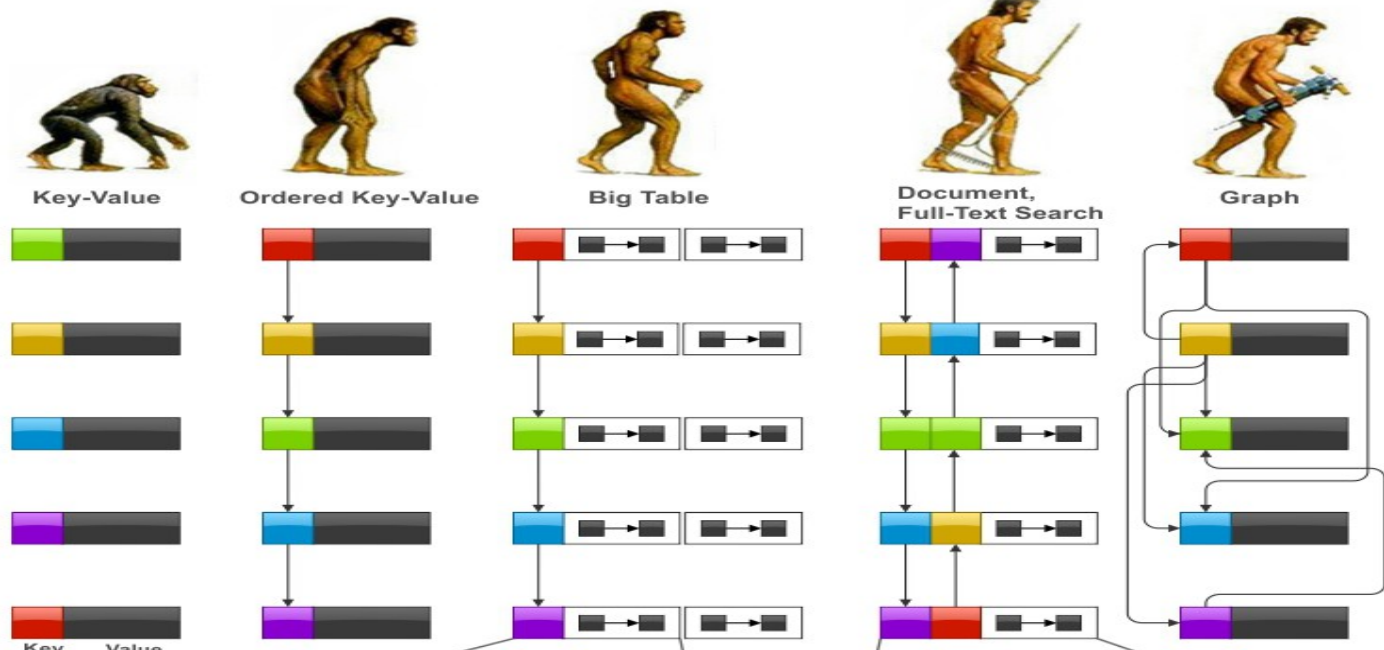
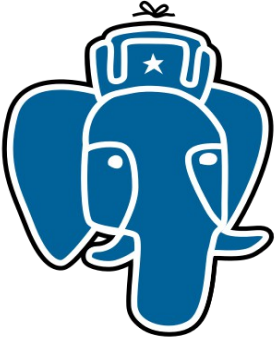
The problem

- The world of data and applications is changing
- BIG DATA (**V**olume of data, **V**elocity of data in-out, **V**ariety of data)
- Web applications are service-oriented
 - Service itself can aggregate data, check consistency of data
 - High concurrency, simple queries
 - Simple database (key-value) is ok
 - Eventual consistency is ok, no ACID overhead
- Application needs faster releases
- NoSQL databases match all of these — scalable, efficient, fault-tolerant, no rigid schema, ready to accept any data.



NoSQL

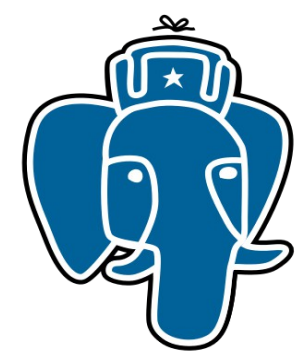
- Key-value databases
 - Ordered k-v for ranges support
- Column family (column-oriented) stores
 - Big Table — value has structure:
 - column families, columns, and timestamped versions (maps-of maps-of maps)
- Document databases
 - Value has arbitrary structure
- Graph databases — evolution of ordered-kv



```
employee" :  
{  
  "name" : "Mohana Pillai"  
  "position" : "Delivery"  
  "projects" : [  
    {  
      "name" : "Easy Signu"  
    }  
  ],  
  "Semi-Structured Data"  
}
```

Plain Text

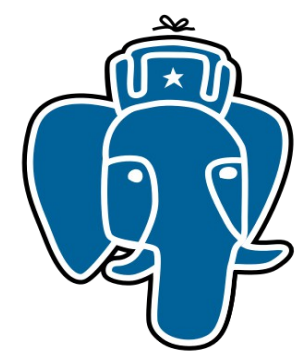
is a confidential word or number
combination used as a code to
identity when accessing
between 8 and 15 characters
number and may not
spaces



The problem

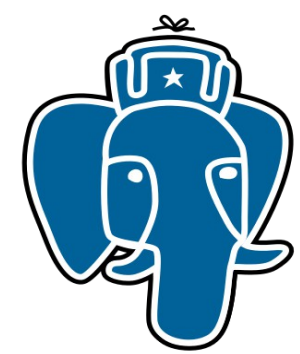
- What if application needs ACID and flexibility of NoSQL ?
- Relational databases work with data with schema known in advance
- It's not easy to change schema online (ALTER TABLE ... ADD COLUMN...)
- Application should wait for schema changing, infrequent releases
- NoSQL uses json format, why not have it in relational database ?

JSON in PostgreSQL
This is the challenge !



Challenge to PostgreSQL !

- Full support of semi-structured data in PostgreSQL
 - Storage
 - Operators and functions
 - Efficiency (fast access to storage, indexes)
 - Integration with CORE (planner, optimiser)
- Actually, PostgreSQL is schema-less database since 2003 — hstore, one of the most popular extension !



Google insights about hstore

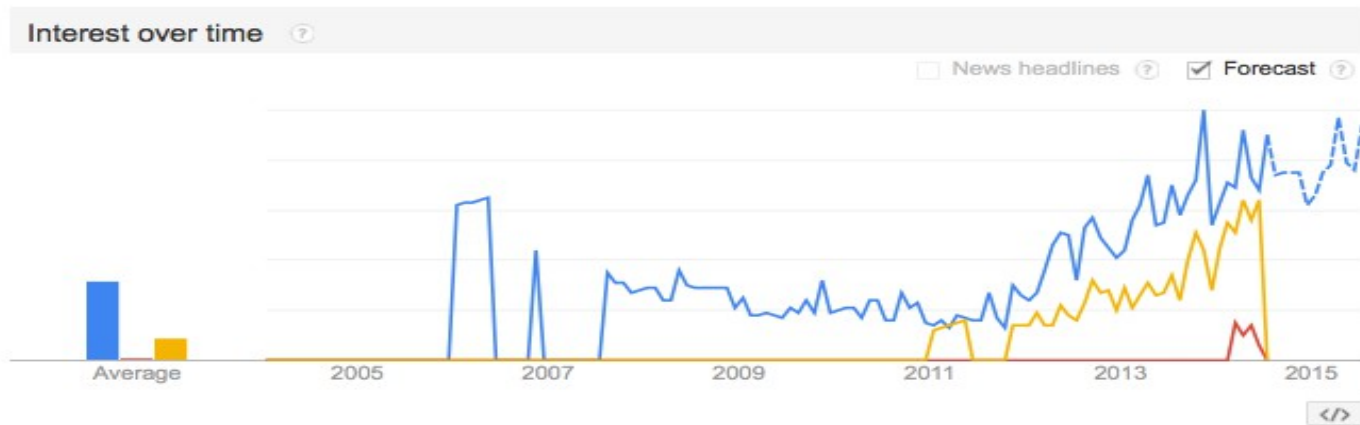
Topics Subscribe <

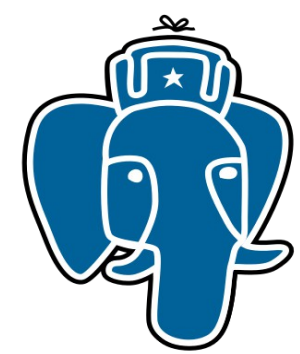
hstore
Search term

jsonb
Search term

json postgresql
Search term

+ Add term

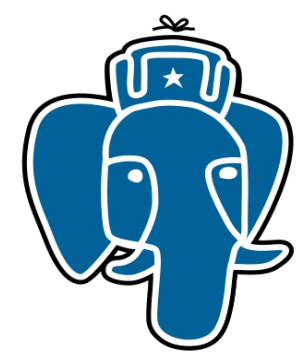




Introduction to Hstore

id	col1	col2	col3	col4	col5	A lot of columns key1, keyN

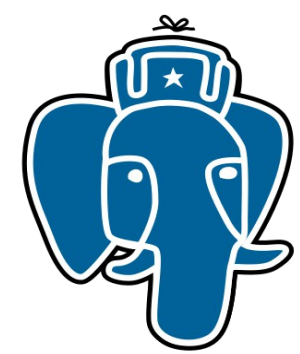
- The problem:
 - Total number of columns may be very large
 - Only several fields are searchable (used in WHERE)
 - Other columns are used only to output
 - These columns may not known in advance
- Solution
 - New data type (hstore), which consists of (key,value) pairs (a'la perl hash)



Introduction to Hstore

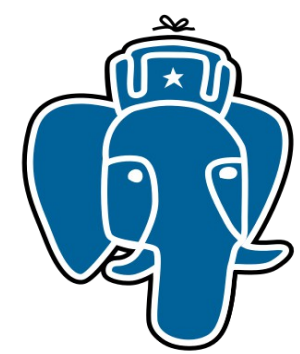
id	col1	col2	col3	col4	col5	Hstore
						key1=>val1, key2=>val2,.....

- Easy to add key=>value pair
- No need change schema, just change hstore.
- Schema-less PostgreSQL in 2003 !



Introduction to hstore

- Hstore — key/value storage (inspired by perl hash)
 - `' a=>1, b=>2 ' ::hstore`
 - Key, value — strings
 - Get value for a key: `hstore -> text`
 - Operators with indexing support (GiST, GIN)
 - Check for key: `hstore ? text`
 - Contains: `hstore @> hstore`
 - [check documentations for more](#)
 - Functions for hstore manipulations (`akeys`, `avals`, `skeys`, `svals`, `each`,.....)
- Hstore provides PostgreSQL schema-less feature !
 - Faster releases, no problem with schema upgrade



History of hstore development

- May 16, 2003 — first version of hstore

```
Date: Fri, 16 May 2003 22:56:14 +0400
From: Teodor Sigaev <teodor@sigaev.ru>
To: Oleg Bartunov <oleg@sai.msu.su>, Alexey Slynko <slynko@tronet.ru>
Cc: E.Rodichev <er@sai.msu.su>
Subject: hash type (hstore)
```

```
Готова первая версия:
zeus:~teodor/hstore.tgz
```

```
README написать не успел, поэтому здесь:
```

```
1 i/o типа hstore
2 операция hstore->text - извлечение значения по ключу text
select 'a=>q, b=>g'->'a';
?
```

```
-----
q
```

```
3 isexists(hstore), isdefined(hstore), delete(hstore,text) - полный перловый аналог
```

```
4 hstore || hstore - конкатенация, аналог в перле %a=( %b, %c );
```

```
5 text=>text - возвращает hstore
```

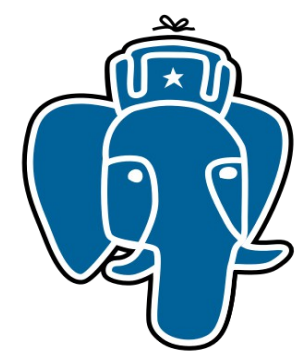
```
select 'a'=>'b';
```

```
?column?
```

```
-----
```

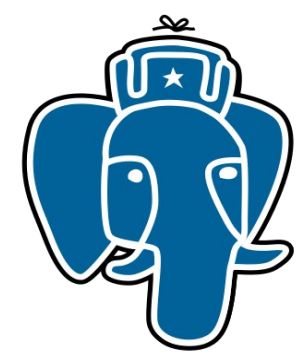
```
"a"=>"b"
```

```
Все примеры есть в sql/hstore.sql
```



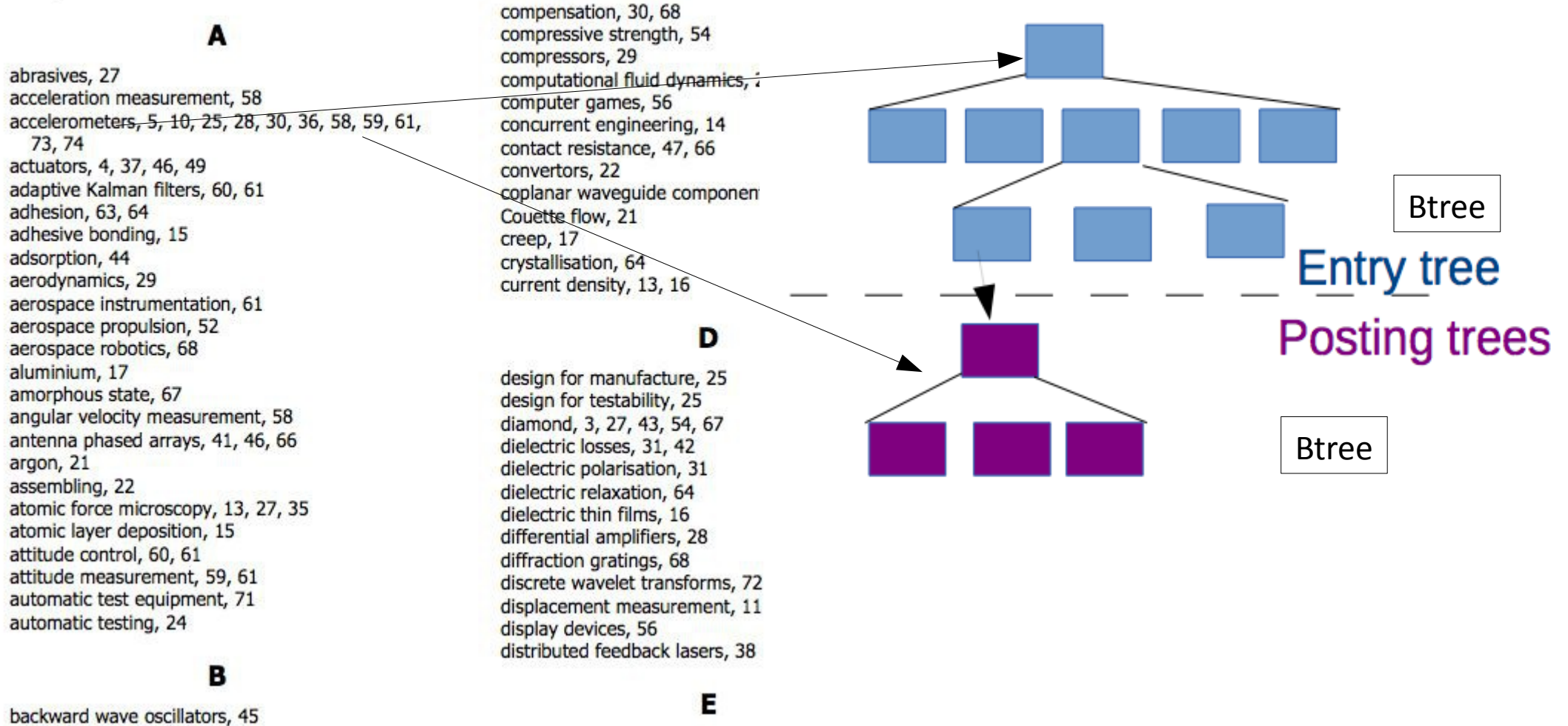
History of hstore development

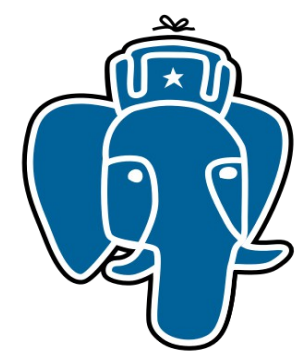
- May 16, 2003 - first (unpublished) version of hstore for PostgreSQL 7.3
- Dec, 05, 2006 - hstore is a part of PostgreSQL 8.2
(thanks, [Hubert Depesz Lubaczewski!](#))
- May 23, 2007 - [GIN index for hstore](#), PostgreSQL 8.3
- Sep, 20, 2010 - Andrew Gierth [improved hstore](#), PostgreSQL 9.0



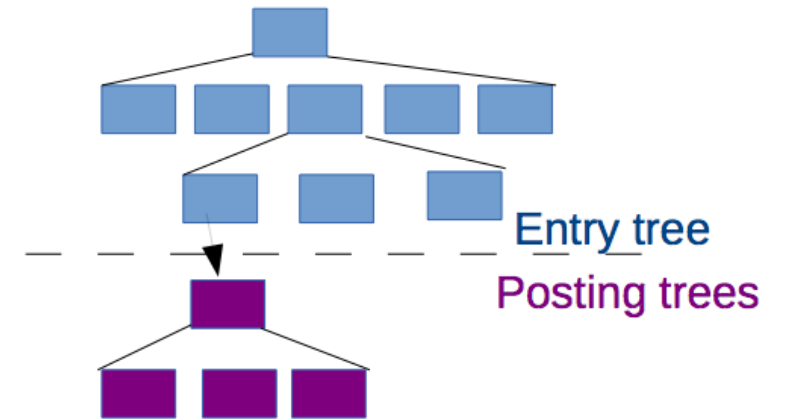
Inverted Index

Report Index





Inverted Index



Report Index

A

abrasives, 27
acceleration measurement, 58
accelerometers, 5, 10, 25, 28, 30, 36, 58, 59, 61, 73, 74
actuators, 4, 37, 46, 49
adaptive Kalman filters, 60, 61
adhesion, 63, 64
adhesive bonding, 15
adsorption, 44
aerodynamics, 29
aerospace instrumentation, 61
aerospace propulsion, 52
aerospace robotics, 68
aluminium, 17
amorphous state, 67
angular velocity measurement
antenna phased arrays, 41, 46
argon, 21
assembling, 22
atomic force microscopy, 13, 21
atomic layer deposition, 15
attitude control, 60, 61
attitude measurement, 59, 61
automatic test equipment, 71
automatic testing, 24

B

backward wave oscillators, 45

compensation, 30, 68
compressive strength, 54
compressors, 29
computational fluid dynamics, 23, 29
computer games, 56
concurrent engineering, 14
contact resistance, 47, 66
convertors, 22
coplanar waveguide components, 40
Couette flow, 21
creep, 17
crystallisation, 64
current density, 13, 16

D

QUERY: compensation accelerometers

INDEX: accelerometers

5,10,25,28,30,36,58,59,61,73,74

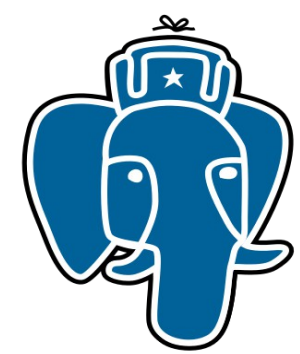
compensation

30,68

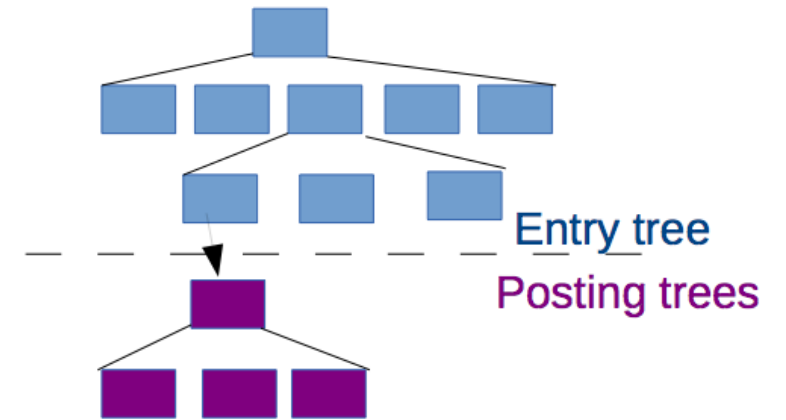
RESULT: **30**

display devices, 30
distributed feedback lasers, 38

E



GIN improvements



- GIN in 9.4 is greatly improved

- Posting lists compression (varbyte encoding) — smaller indexes

- 9.3: always 6 bytes (4 bytes blockNumber , 2 bytes offset): **90 bytes**

(0,8) (0,14) (0,17) (0,22) (0,26) (0,33) (0,34) (0,35) (0,45) (0,47) (0,48) (1,3) (1,4)
(1,6) (1,8)

- 9.4: 1-6 bytes per each item, deltas from previous item: **21 bytes**

(0,8) +6 +3 +5 +4 +7 +1 +1 +10 +2 +1 +2051 +1+2 +2

SELECT g % 10 FROM generate_series(1,10000000) g; **11Mb vs 58Mb**

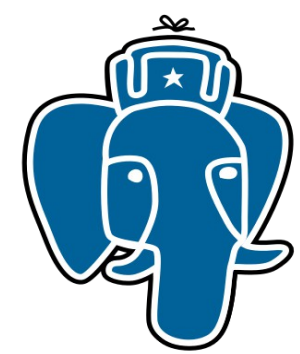
- Fast scan of posting lists - «rare & frequent» queries much faster

- 9.3: read posting lists for «rare» and «frequent» and join them

Time(frequent & rare) ~ Time(frequent)

- 9.4: start from posting list for «rare» and skip «frequent» list if no match

Time(frequent & rare) ~ Time(rare)

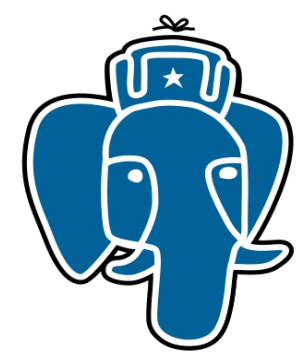


Hstore is DEAD ? No !

- How hstore benefits by GIN improvement in 9.4 ?

GIN stands for Generalized Inverted Index, so virtually all data types, which use GIN, get benefit !

- Default hstore GIN opclass considers keys and values separately
- Keys are «frequent», value are «rare»
- Contains query: `hstore @> 'key=>value'` improved a lot for «rare» values
- Index size is smaller, less io



Hstore 9.3 vs 9.4

Total: 7240858 geo records:

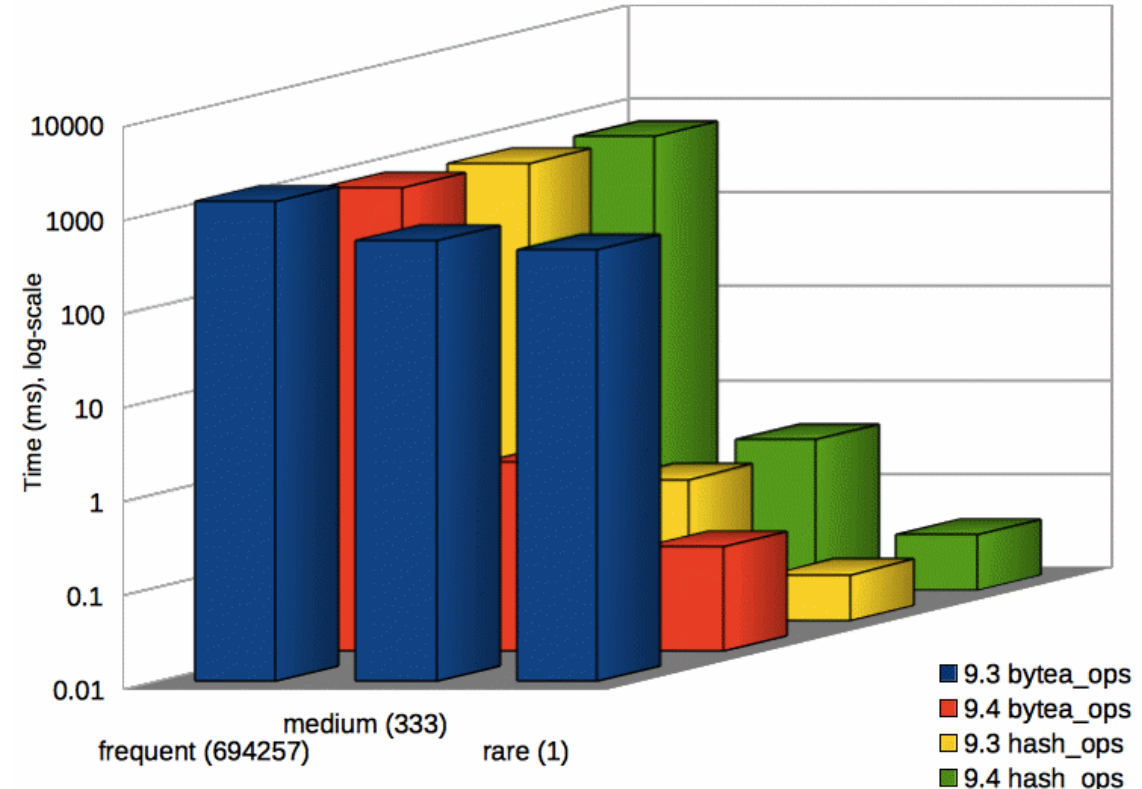
```
"fcode"=>"RFSU",  
"point"=>"(8.85,112.53333)",  
"fclass"=>"U",  
"asciiname"=>"London Reefs",  
"elevation"=>NULL,  
"geonameid"=>"1879967",  
"population"=>"0"
```

Query:

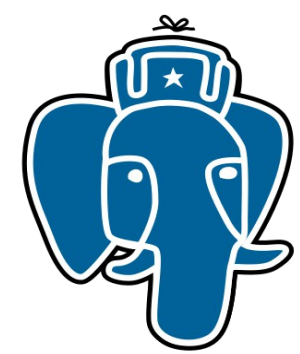
```
SELECT count(*) FROM geo  
WHERE geo @> 'fcode=>STM';
```

opclass	frequent (694257)	medium (333)	rare (1)
9.3 bytea_ops	1353.844	511.196	402.662
9.4 bytea_ops	878.875	1.031	0.13
9.3 hash_ops	755.458	0.321	0.031
9.4 hash_ops	687.626	0.4	0.039

Hstore GIN opclass performance: 9.3 vs 9.4



gin_hstore_ops: index keys and values
gin_hstore_bytea_ops = gin_hstore_ops, no collation comparison
gin_hstore_hash_ops: index hash(key.value)



Hstore 9.3 vs 9.4

9.3

Name	Type	Owner	Table	Size
geo	table	postgres		1352 MB
geo_hstore_bytea_ops	index	postgres	geo	1680 MB
geo_hstore_hash_ops_idx	index	postgres	geo	1073 MB

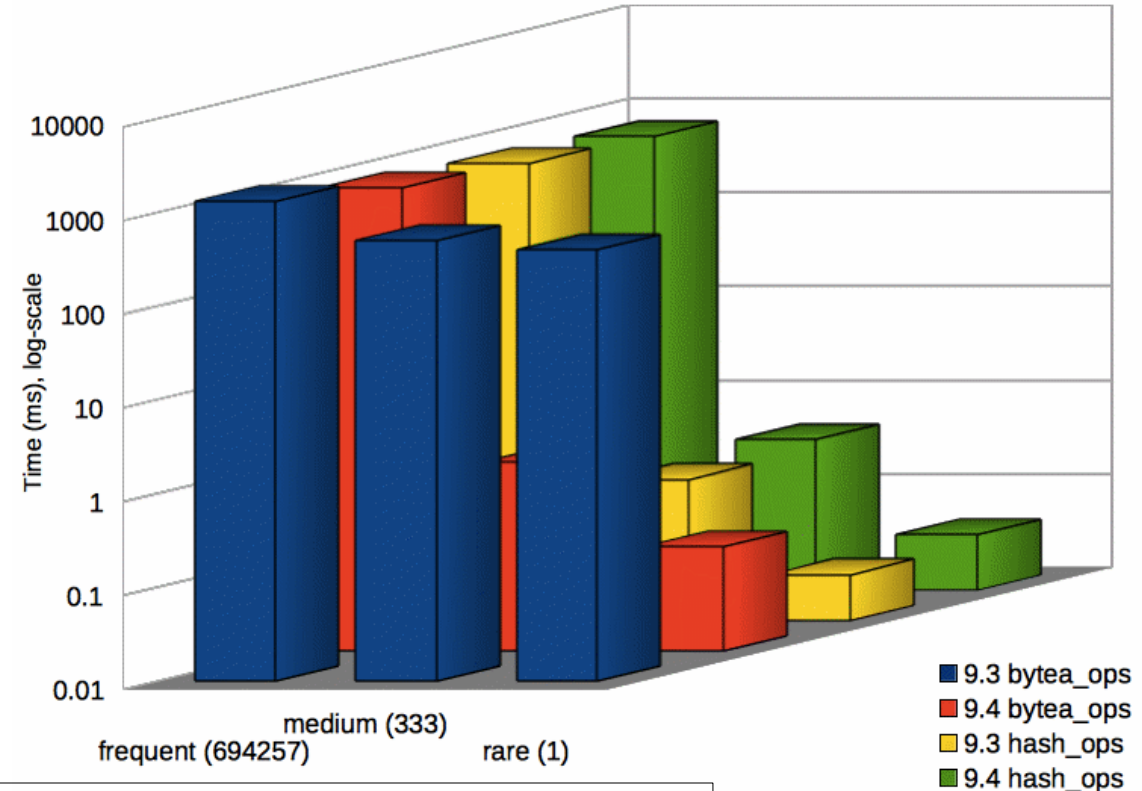
9.4

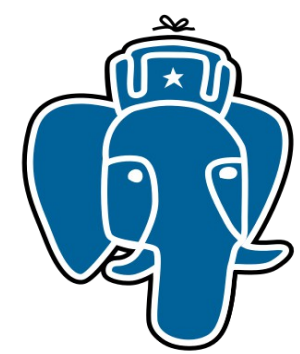
Name	Type	Owner	Table	Size
geo	table	postgres		1352 MB
geo_hstore_bytea_ops	index	postgres	geo	1296 MB
geo_hstore_hash_ops_idx	index	postgres	geo	925 MB

```
CREATE OPERATOR CLASS gin_hstore_bytea_ops FOR TYPE hstore
.....
FUNCTION 1 byteacmp(bytea,bytea),
.....
STORAGE bytea;
CREATE INDEX: 239 s Much faster comparison (no collation)
```

```
CREATE OPERATOR CLASS gin_hstore_ops FOR TYPE hstore
.....
FUNCTION 1 btttextcmp(text,text),,
.....
STORAGE text;
CREATE INDEX: 2870 s
```

Hstore GIN oclass performance: 9.3 vs 9.4





Hstore 9.3 vs 9.4

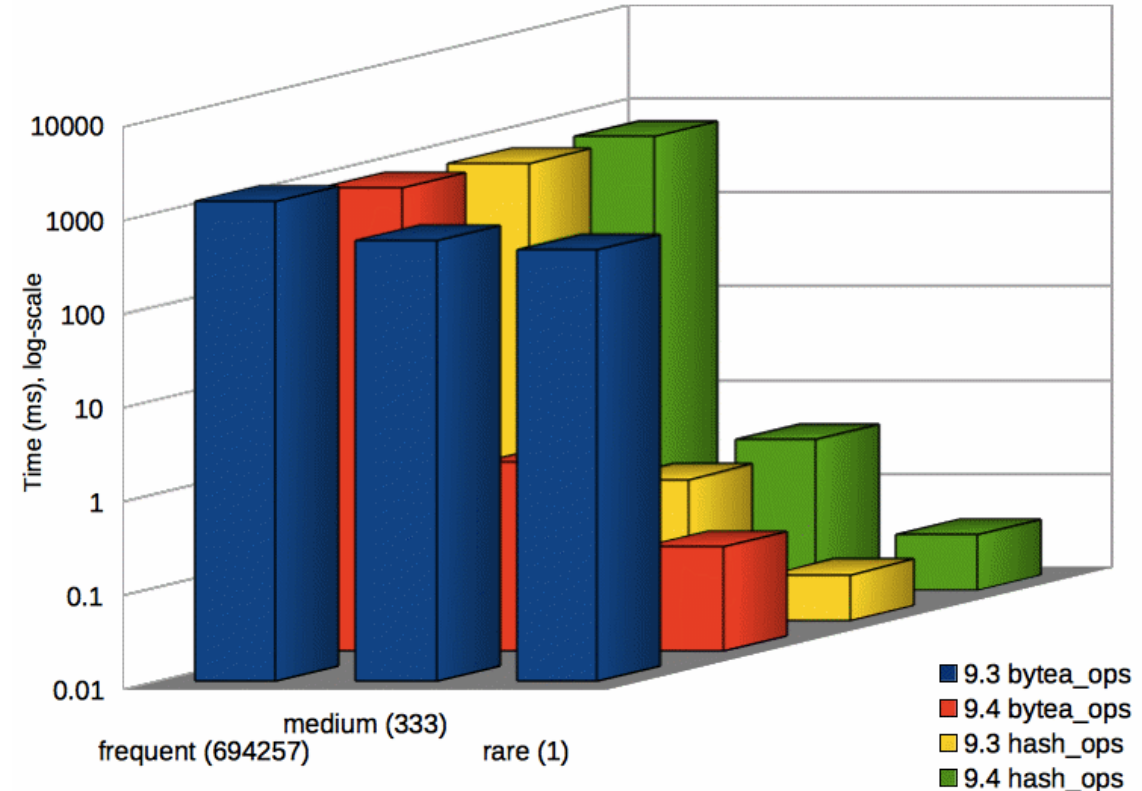
SUMMARY:

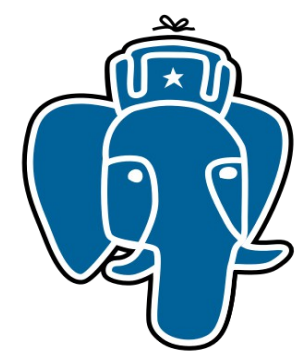
- 9.4 GIN posting list compression: indexes are smaller
- 9.4 GIN is smart regarding 'freq & rare' queries: time (freq & rare) ~ time (rare) instead of time (freq & rare) ~ time (freq)
- gin_hstore_hash_ops is good on 9.3 & 9.4 and faster default gin opclass
- Use gin_hstore_bytea_ops instead of default gin_hstore_ops — much faster create index

Get hstore_ops from:

from https://github.com/akorotkov/hstore_ops

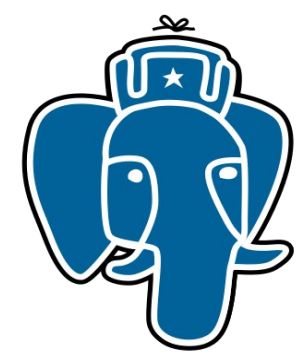
Hstore GIN opclass performance: 9.3 vs 9.4





Introduction to hstore

- Hstore benefits
 - In provides a flexible model for storing a semi-structured data in relational database
 - hstore has binary storage and rich set of operators and functions, indexes
- Hstore drawbacks
 - Too simple model !
Hstore key-value model doesn't supports tree-like structures as json (introduced in 2006, 3 years after hstore)
- Json — popular and standartized (ECMA-404 The JSON Data Interchange Standard, JSON RFC-7159)
- Json — PostgreSQL 9.2, textual storage



Hstore vs Json

- hstore is faster than json even on simple data

```
CREATE TABLE hstore_test AS (SELECT  
'a=>1, b=>2, c=>3, d=>4, e=>5'::hstore AS v  
FROM generate_series(1,1000000));
```

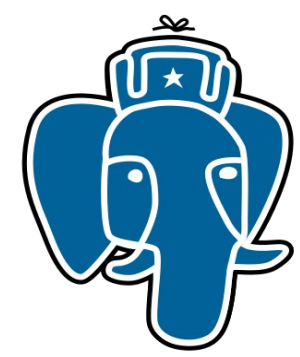
```
CREATE TABLE json_test AS (SELECT  
'{"a":1, "b":2, "c":3, "d":4, "e":5}'::json AS v  
FROM generate_series(1,1000000));
```

```
SELECT sum((v->'a')::text::int) FROM json_test;
```

851.012 ms

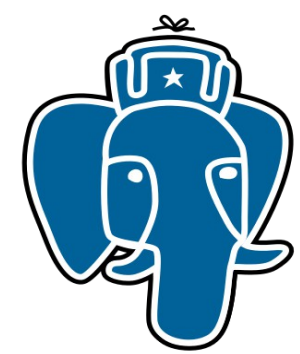
```
SELECT sum((v->'a')::int) FROM hstore_test;
```

330.027 ms



Hstore vs Json

- PostgreSQL already has json since 9.2, which supports document-based model, but
 - It's slow, since it has no binary representation and needs to be parsed every time
 - Hstore is fast, thanks to binary representation and index support
 - It's possible to convert hstore to json and vice versa, but current hstore is limited to key-value
 - **Need hstore with document-based model. Share it's binary representation with json !**



Nested hstore

abstract ▾



Oleg Bartunov <obartunov@gmail.com>

12/18/12 ☆



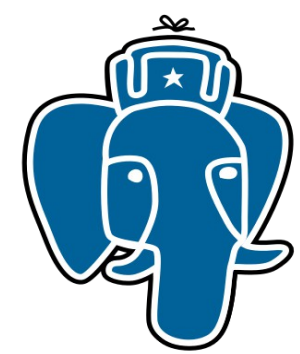
to Teodor ▾

Поправь, дополни.

Title: One step forward true json data type. **Nested hstore** with array support.

We present a prototype of **nested hstore** data type with array support. We consider the new **hstore** as a step forward true json data type.

Recently, PostgreSQL got json data type, which basically is a string storage with validity checking for stored values and some related functions. To be a real data type, it has to have a binary representation, which could be a big project if started from scratch. **Hstore** is a popular data type, we developed years ago to facilitate working with semi-structured data in PostgreSQL. Our idea is to extend **hstore** to be **nested** (value can be **hstore**) data type and add support of arrays, so its binary representation can be shared with json. We present a working prototype of a new **hstore** data type and discuss some design and implementation issues.



Nested hstore & jsonb

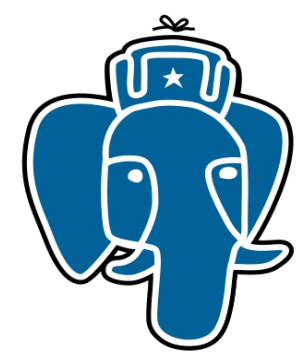
- Nested hstore at PGCon-2013, Ottawa, Canada (May 24) — thanks Engine Yard for support !

One step forward true json data type. Nested hstore with arrays support

- Binary storage for nested data at PGCon Europe — 2013, Dublin, Ireland (Oct 29)

Binary storage for nested data structures and application to hstore data type

- November, 2013 — binary storage was reworked, nested hstore and jsonb share the same storage. Andrew Dunstan joined the project.
- January, 2014 - binary storage moved to core

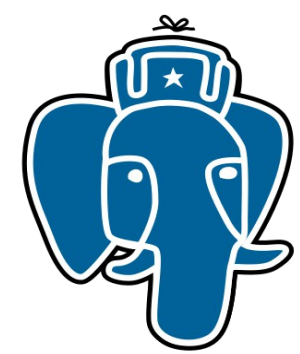


Nested hstore & jsonb

- Feb-Mar, 2014 - Peter Geoghegan joined the project, nested hstore was cancelled in favour to jsonb ([Nested hstore patch for 9.3](#)).
- Mar 23, 2014 Andrew Dunstan committed jsonb to 9.4 branch !
[pgsql: Introduce jsonb, a structured format for storing json.](#)

Introduce jsonb, a structured format for storing json.

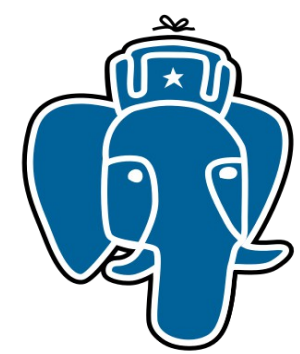
The new format accepts exactly the same data as the json type. However, it is stored in a format that does not require reparsing the original text in order to process it, making it much more suitable for indexing and other operations. Insignificant whitespace is discarded, and the order of object keys is not preserved. Neither are duplicate object keys kept - the later value for a given key is the only one stored.



Jsonb vs Json

```
SELECT '{"c":0, "a":2,"a":1}'::json, '{"c":0, "a":2,"a":1}'::jsonb;
      json                |      jsonb
-----+-----
{"c":0, "a":2,"a":1} | {"a": 1, "c": 0}
(1 row)
```

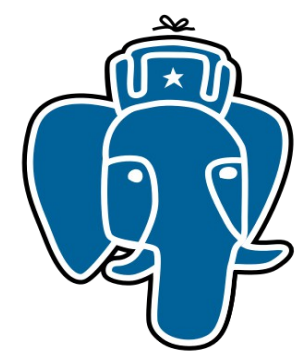
- json: textual storage «as is»
- jsonb: no whitespaces
- jsonb: no duplicate keys, last key win
- jsonb: keys are sorted



Jsonb vs Json

- Data
 - 1,252,973 Delicious bookmarks
- Server
 - MBA, 8 GB RAM, 256 GB SSD
- Test
 - Input performance - copy data to table
 - Access performance - get value by key
 - Search performance contains @> operator

```
{
  "author": "mcasas1",
  "comments": "http://delicious.com/url/b5b3cbf9a9176fe43c27d7b4af94a422",
  "guidislink": false,
  "id": "http://delicious.com/url/b5b3cbf9a9176fe43c27d7b4af94a422#mcasas1",
  "link": "http://www.theatermania.com/broadway/",
  "links": [
    {
      "href": "http://www.theatermania.com/broadway/",
      "rel": "alternate",
      "type": "text/html"
    }
  ],
  "source": {},
  "tags": [
    {
      "label": null,
      "scheme": "http://delicious.com/mcasas1/",
      "term": "NYC"
    }
  ],
  "title": "TheaterMania",|
  "title_detail": {
    "base": "http://feeds.delicious.com/v2/rss/recent?min=1&count=100",
    "language": null,
    "type": "text/plain",
    "value": "TheaterMania"
  },
  "updated": "Tue, 08 Sep 2009 23:28:55 +0000",
  "wfw_commentrss": "http://feeds.delicious.com/v2/rss/url/b5b3cbf9a9176fe43c27d7b4af94a422"
}
```



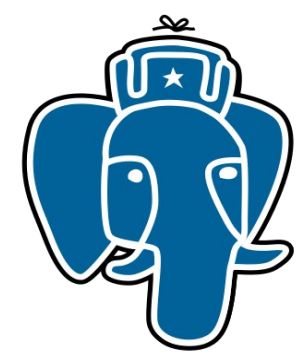
Jsonb vs Json

- Data
 - 1,252,973 bookmarks from Delicious in json format (js)
 - The same bookmarks in jsonb format (jb)
 - The same bookmarks as text (tx)

```
=# \dt+
```

```
List of relations
```

Schema	Name	Type	Owner	Size	Description
public	jb	table	postgres	1374 MB	overhead is < 4%
public	js	table	postgres	1322 MB	
public	tx	table	postgres	1322 MB	



Jsonb vs Json

- Input performance (parser)

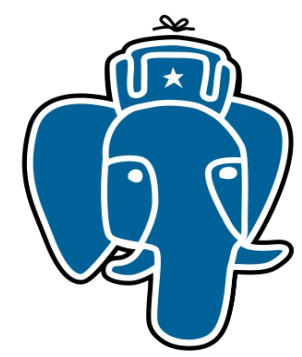
Copy data (1,252,973 rows) as text, json,jsonb

copy tt from '/path/to/test.dump'

Text: 34 s - as is

Json: 37 s - json validation

Jsonb: 43 s - json validation, binary storage



Jsonb vs Json (binary storage)

- Access performance — get value by key

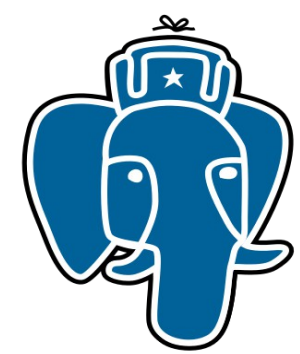
- Base: `SELECT js FROM js;`
- Jsonb: `SELECT j->>'updated' FROM jb;`
- Json: `SELECT j->>'updated' FROM js;`

Base: 0.6 s

Jsonb: 1 s 0.4

Json: 9.6 s 9

Jsonb ~ 20X faster Json

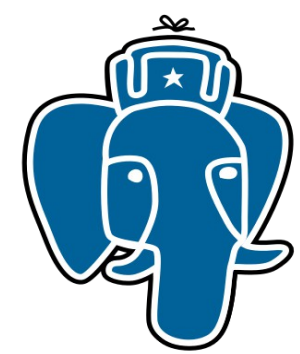


Jsonb vs Json

```
EXPLAIN ANALYZE SELECT count(*) FROM js WHERE js #>>' {tags,0,term}' = 'NYC';  
QUERY PLAN
```

```
-----  
Aggregate (cost=187812.38..187812.39 rows=1 width=0)  
(actual time=10054.602..10054.602 rows=1 loops=1)  
  -> Seq Scan on js (cost=0.00..187796.88 rows=6201 width=0)  
(actual time=0.030..10054.426 rows=123 loops=1)  
    Filter: ((js #>> '{tags,0,term}'::text[]) = 'NYC'::text)  
    Rows Removed by Filter: 1252850  
Planning time: 0.078 ms  
Execution runtime: 10054.635 ms  
(6 rows)
```

**Json: no contains @> operator,
search first array element**

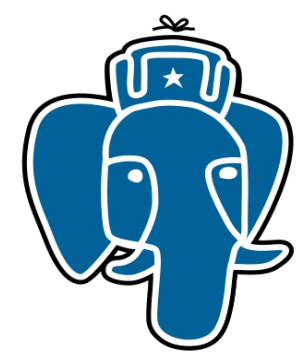


Jsonb vs Json (binary storage)

```
EXPLAIN ANALYZE SELECT count(*) FROM jb WHERE jb @> '{"tags":[{"term":"NYC"}]}'::jsonb;  
QUERY PLAN
```

```
-----  
Aggregate (cost=191521.30..191521.31 rows=1 width=0)  
(actual time=1263.201..1263.201 rows=1 loops=1)  
  -> Seq Scan on jb (cost=0.00..191518.16 rows=1253 width=0)  
    (actual time=0.007..1263.065 rows=285 loops=1)  
      Filter: (jb @> '{"tags": [{"term": "NYC"}]}'::jsonb)  
      Rows Removed by Filter: 1252688  
    Planning time: 0.065 ms  
    Execution runtime: 1263.225 ms  
    Execution runtime: 10054.635 ms  
(6 rows)
```

Jsonb ~ 10X faster Json



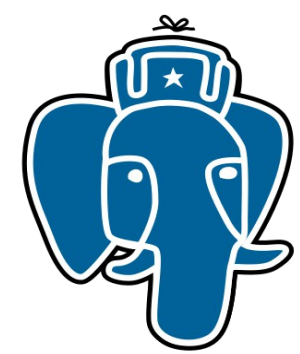
Jsonb vs Json (GIN: key && value)

```
CREATE INDEX gin_jb_idx ON jb USING gin(jb);
```

```
EXPLAIN ANALYZE SELECT count(*) FROM jb WHERE jb @> '{"tags":[{"term":"NYC"}]}'::jsonb;  
QUERY PLAN
```

```
-----  
Aggregate (cost=4772.72..4772.73 rows=1 width=0)  
(actual time=8.486..8.486 rows=1 loops=1)  
  -> Bitmap Heap Scan on jb (cost=73.71..4769.59 rows=1253 width=0)  
(actual time=8.049..8.462 rows=285 loops=1)  
    Recheck Cond: (jb @> '{"tags": [{"term": "NYC"}]}'::jsonb)  
    Heap Blocks: exact=285  
    -> Bitmap Index Scan on gin_jb_idx (cost=0.00..73.40 rows=1253 width=0)  
(actual time=8.014..8.014 rows=285 loops=1)  
      Index Cond: (jb @> '{"tags": [{"term": "NYC"}]}'::jsonb)  
Planning time: 0.115 ms  
Execution runtime: 8.515 ms           Execution runtime: 10054.635 ms  
(8 rows)
```

Jsonb ~ 150X faster Json



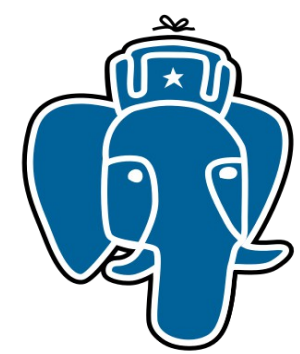
Jsonb vs Json (GIN: hash path.value)

```
CREATE INDEX gin_jb_path_idx ON jb USING gin(jb jsonb_path_ops);
```

```
EXPLAIN ANALYZE SELECT count(*) FROM jb WHERE jb @> '{"tags":[{"term":"NYC"}]}'::jsonb;  
QUERY PLAN
```

```
-----  
Aggregate (cost=4732.72..4732.73 rows=1 width=0)  
(actual time=0.644..0.644 rows=1 loops=1)  
  -> Bitmap Heap Scan on jb (cost=33.71..4729.59 rows=1253 width=0)  
(actual time=0.102..0.620 rows=285 loops=1)  
    Recheck Cond: (jb @> '{"tags": [{"term": "NYC"}]}'::jsonb)  
    Heap Blocks: exact=285  
    -> Bitmap Index Scan on gin_jb_path_idx  
(cost=0.00..33.40 rows=1253 width=0) (actual time=0.062..0.062 rows=285 loops=1)  
      Index Cond: (jb @> '{"tags": [{"term": "NYC"}]}'::jsonb)  
Planning time: 0.056 ms  
Execution runtime: 0.668 ms          Execution runtime: 10054.635 ms  
(8 rows)
```

Jsonb ~ 1800X faster Json



MongoDB 2.6.0

- Load data - ~13 min **SLOW !**

Jsonb 43 s

```
mongoimport --host localhost -c js --type json < delicious-rss-1250k
```

```
2014-04-08T22:47:10.014+0400
```

```
3700
```

```
1233/second
```

```
...
```

```
2014-04-08T23:00:36.050+0400
```

```
1252000
```

```
1547/second
```

```
2014-04-08T23:00:36.565+0400 check 9 1252973
```

```
2014-04-08T23:00:36.566+0400 imported 1252973 objects
```

- Search - ~ 1s (seqscan) **THE SAME**

```
db.js.find({tags: {$elemMatch:{ term: "NYC"}}}).count()
```

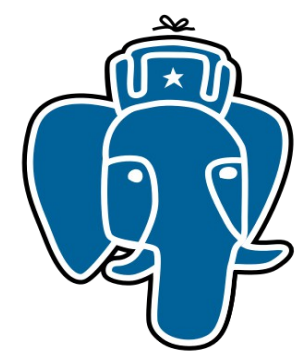
```
285
```

```
-- 980 ms
```

- Search - ~ 1ms (indexscan) **Jsonb 0.7ms**

```
db.js.ensureIndex( {"tags.term" : 1} )
```

```
db.js.find({tags: {$elemMatch:{ term: "NYC"}}}).
```



Summary: PostgreSQL 9.4 vs Mongo 2.6.0

- Operator contains @>

- json : 10 s seqscan
- jsonb : 8.5 ms GIN jsonb_ops
- **jsonb : 0.7 ms GIN jsonb_path_ops**
- mongo : 1.0 ms btree index

- Index size

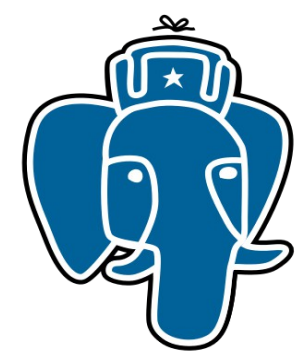
- jsonb_ops - 636 Mb (no compression, 815Mb)
- jsonb_path_ops - 295 Mb
- jsonb_path_ops (tags) - 44 Mb USING gin((jb->'tags') jsonb_path_ops)
- mongo (tags) - 387 Mb
- mongo (tags.term) - 100 Mb

- Table size

- postgres : 1.3Gb
- mongo : 1.8Gb

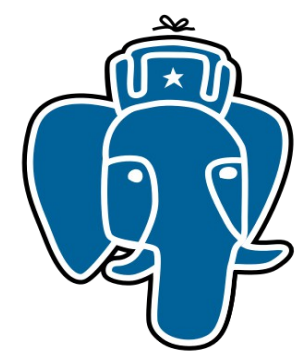
- Input performance:

- Text : 34 s
- Json : 37 s
- Jsonb : 43 s
- mongo : 13 m



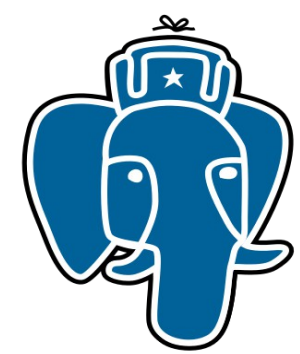
Jsonb (Apr, 2014)

- Documentation
 - [JSON Types, JSON Functions and Operators](#)
- There are many functionality left in nested hstore
 - Can be an extension
- Need query language for jsonb
 - `<, >, && ...` operators for values
 - `a.b.c.d && [1,2,10]`
 - Structural queries on paths
 - `*.d && [1,2,10]`
 - Indexes !



Jsonb query

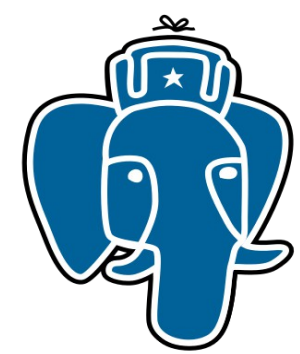
- Currently, one can search jsonb data using
 - Contains operators - `jsonb @> jsonb`, `jsonb <@ jsonb` (GIN indexes)
`jb @> '{"tags": [{"term": "NYC"}]}'::jsonb`
Keys should be specified from root
 - Equivalence operator — `jsonb = jsonb` (GIN indexes)
 - Exists operators — `jsonb ? text`, `jsonb ?! text[]`, `jsonb ?& text[]` (GIN indexes)
`jb WHERE jb ?| '{tags,links}'`
Only root keys supported
 - Operators on jsonb parts (functional indexes)
`SELECT ('{"a": {"b":5}}'::jsonb -> 'a'->>'b')::int > 2;`
`CREATE INDEXUSING BTREE ((jb->'a'->>'b')::int);`
Very cumbersome, too many functional indexes



Jsonb query

- Need Jsonb query language
 - More operators on keys, values
 - Types support
 - Schema support (constraints on keys, values)
 - Indexes support
- Introduce Jsquery - textual data type and @@ match operator

jsonb @@ jsquery



Jsonb query language (Jsquery)

```
Expr ::= path value_expr
      | path HINT value_expr
      | NOT expr
      | NOT HINT value_expr
      | NOT value_expr
      | path '(' expr ')
      | '(' expr ')
      | expr AND expr
      | expr OR expr
```

```
value_expr
  ::= '=' scalar_value
     | IN '(' value_list ')'
     | '=' array
     | '=' '*'
     | '<' NUMERIC
     | '<' '=' NUMERIC
     | '>' NUMERIC
     | '>' '=' NUMERIC
     | '@' '>' array
     | '<' '@' array
     | '&' '&' array
     | IS ARRAY
     | IS NUMERIC
     | IS OBJECT
     | IS STRING
     | IS BOOLEAN
```

```
path ::= key
      | path '.' key_any
      | NOT '.' key_any

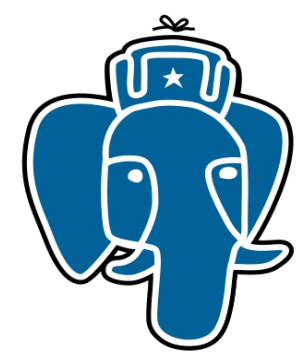
key   ::= '*'
      | '#'
      | '%'
      | '$'
      | STRING
      | .....

key_any ::= key
        | NOT
```

```
value_list
  ::= scalar_value
     | value_list ',' scalar_value

array ::= '[' value_list ']'

scalar_value
  ::= null
     | STRING
     | true
     | false
     | NUMERIC
     | OBJECT
     | .....
```



Jsonb query language (Jsquery)

- # - any element array

```
SELECT '{"a": {"b": [1,2,3]}}'::jsonb @@ 'a.b.# = 2';
```

- % - any key

```
SELECT '{"a": {"b": [1,2,3]}}'::jsonb @@ '%.b.# = 2';
```

- * - anything

```
SELECT '{"a": {"b": [1,2,3]}}'::jsonb @@ '*.# = 2';
```

- \$ - current element

```
select '{"a": {"b": [1,2,3]}}'::jsonb @@ 'a.b.# ($ = 2 OR $ < 3)';
```

- Use "double quotes" for key !

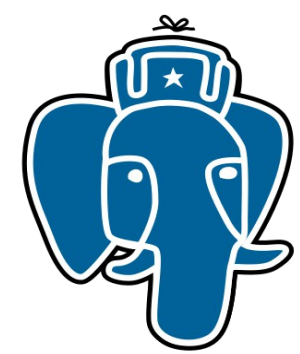
```
select 'a1."12222" < 111'::jsquery;
```

```
path ::= key  
      | path '.' key_any  
      | NOT '.' key_any
```

```
key ::= '*'  
     | '#'  
     | '%'  
     | '$'  
     | STRING
```

.....

```
key_any ::= key  
         | NOT
```



Jsonb query language (Jsquery)

- Scalar

```
select '{"a": {"b": [1,2,3]}}'::jsonb @@ 'a.b.# IN (1,2,5)';
```

- Test for key existence

```
select '{"a": {"b": [1,2,3]}}'::jsonb @@ 'a.b = *';
```

- Array overlap

```
select '{"a": {"b": [1,2,3]}}'::jsonb @@ 'a.b && [1,2,5]';
```

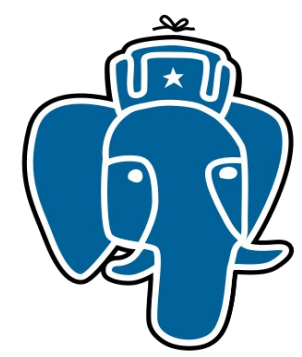
- Array contains

```
select '{"a": {"b": [1,2,3]}}'::jsonb @@ 'a.b @> [1,2]';
```

- Array contained

```
select '{"a": {"b": [1,2,3]}}'::jsonb @@ 'a.b <@ [1,2,3,4,5]';
```

```
value_expr
 ::= '=' scalar_value
 | IN '(' value_list ')'
 | '=' array
 | '=' '*'
 | '<' NUMERIC
 | '<' '=' NUMERIC
 | '>' NUMERIC
 | '>' '=' NUMERIC
 | '@' '>' array
 | '<' '@' array
 | '&' '&' array
 | IS ARRAY
 | IS NUMERIC
 | IS OBJECT
 | IS STRING
 | IS BOOLEAN
```



Jsonb query language (Jsqquery)

- Type checking

```
select '{"x": true}' @@ 'x IS boolean'::jsquery,  
       '{"x": 0.1}'  @@ 'x IS numeric'::jsquery;  
?column? | ?column?  
-----+-----  
t         | t
```

```
select '{"a":{"a":1}}' @@ 'a IS object'::jsquery;  
?column?  
-----  
t
```

```
select '{"a":["xxx"]}' @@ 'a IS array'::jsquery, '["xxx"]' @@ '$ IS array'::jsquery;  
?column? | ?column?  
-----+-----  
t         | t
```

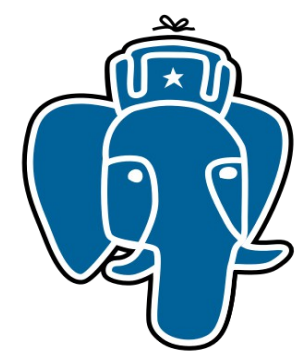
IS BOOLEAN

IS NUMERIC

IS ARRAY

IS OBJECT

IS STRING



Jsonb query language (Jsqquery)

- How many products are similar to "B000089778" and have product_sales_rank in range between 10000-20000 ?

- SQL

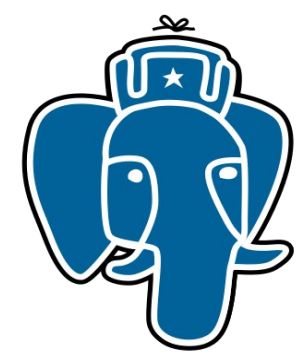
```
SELECT count(*) FROM jr WHERE (jr->>'product_sales_rank')::int > 10000  
and (jr->> 'product_sales_rank')::int < 20000 and  
....boring stuff
```

- Jsqquery

```
SELECT count(*) FROM jr WHERE jr @@ 'similar_product_ids &&  
["B000089778"] AND product_sales_rank( $ > 10000 AND $ < 20000)'
```

- MongoDB

```
db.reviews.find( { $and :[ {similar_product_ids: { $in ["B000089778"]}},  
{product_sales_rank:{$gt:10000, $lt:20000}}] } ).count()
```



Jsonb query language (Jsquery)

```
explain( analyze, buffers) select count(*) from jb where jb @> '{"tags": [{"term": "NYC"}]}'::jsonb;  
QUERY PLAN
```

Aggregate (cost=191517.30..191517.31 rows=1 width=0) (actual time=1039.422..1039.423 rows=1 loops=1)

Buffers: shared hit=97841 read=78011

-> Seq Scan on jb (cost=0.00..191514.16 rows=1253 width=0) (actual time=0.006..1039.310 rows=285 loops=1)

Filter: (jb @> '{"tags": [{"term": "NYC"}]}'::jsonb)

Rows Removed by Filter: 1252688

Buffers: shared hit=97841 read=78011

Planning time: 0.074 ms

Execution time: 1039.444 ms

```
explain( analyze, costs off) select count(*) from jb where jb @@ 'tags.#.term = "NYC";  
QUERY PLAN
```

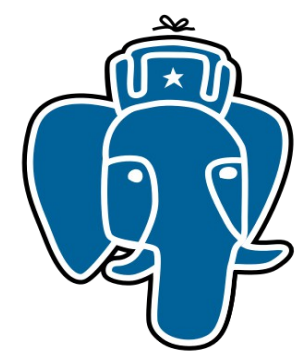
Aggregate (actual time=891.707..891.707 rows=1 loops=1)

-> Seq Scan on jb (actual time=0.010..891.553 rows=285 loops=1)

Filter: (jb @@ "tags".#.term = "NYC"::jsquery)

Rows Removed by Filter: 1252688

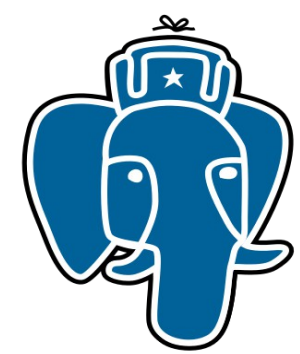
Execution time: 891.745 ms



Jsquery (indexes)

- GIN opclasses with jsquery support
 - jsonb_value_path_ops — use Bloom filtering for key matching
`{"a":{"b":{"c":10}}}` → 10.(bloom(a) or bloom(b) or bloom(c))
 - Good for key matching (wildcard support) , not good for range query
 - jsonb_path_value_ops — hash path (like jsonb_path_ops)
`{"a":{"b":{"c":10}}}` → hash(a.b.c).10
 - No wildcard support, no problem with ranges

Schema	Name	Type	Owner	Table	Size	Description
public	jb	table	postgres		1374 MB	
public	jb_value_path_idx	index	postgres	jb	306 MB	
public	jb_gin_idx	index	postgres	jb	544 MB	
public	jb_path_value_idx	index	postgres	jb	306 MB	
public	jb_path_idx	index	postgres	jb	251 MB	



Jsquery (indexes)

```
explain( analyze, costs off) select count(*) from jb where jb @@ 'tags.#.term = "NYC";
```

QUERY PLAN

Aggregate (actual time=0.609..0.609 rows=1 loops=1)

-> Bitmap Heap Scan on jb (actual time=0.115..0.580 rows=285 loops=1)

Recheck Cond: (jb @@ "'tags'.#.'term' = 'NYC'::jsquery)

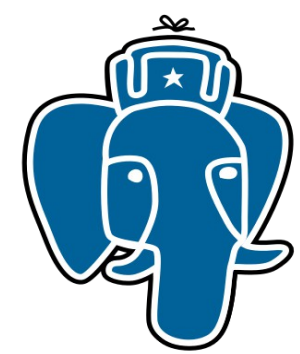
Heap Blocks: exact=285

-> Bitmap Index Scan on jb_value_path_idx (actual time=0.073..0.073 rows=285 loops=1)

Index Cond: (jb @@ "'tags'.#.'term' = 'NYC'::jsquery)

Execution time: 0.634 ms

(7 rows)



Jsquery (indexes)

```
explain( analyze, costs off) select count(*) from jb where jb @@ '*.term = "NYC";
```

QUERY PLAN

Aggregate (actual time=0.688..0.688 rows=1 loops=1)

-> Bitmap Heap Scan on jb (actual time=0.145..0.660 rows=285 loops=1)

Recheck Cond: (jb @@ '*."term" = "NYC"::jsquery)

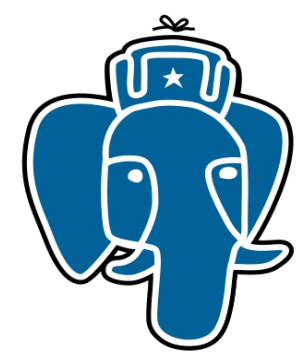
Heap Blocks: exact=285

-> Bitmap Index Scan on jb_value_path_idx (actual time=0.113..0.113 rows=285 loops=1)

Index Cond: (jb @@ '*."term" = "NYC"::jsquery)

Execution time: 0.716 ms

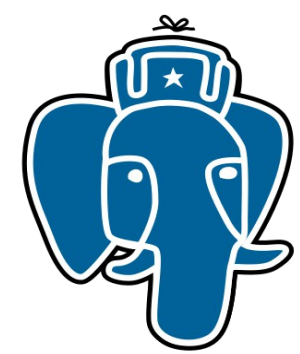
(7 rows)



Citus dataset

- 3023162 reviews from Citus 1998-2000 years
- 1573 MB

```
{
  "customer_id": "AE22YDHSBFYIP",
  "product_category": "Business & Investing",
  "product_group": "Book",
  "product_id": "1551803542",
  "product_sales_rank": 11611,
  "product_subcategory": "General",
  "product_title": "Start and Run a Coffee Bar (Start & Run a)",
  "review_date": {
    "$date": 31363200000
  },
  "review_helpful_votes": 0,
  "review_rating": 5,
  "review_votes": 10,
  "similar_product_ids": [
    "0471136174",
    "0910627312",
    "047112138X",
    "0786883561",
    "0201570483"
  ]
}
```



Jsquery (indexes)

```
explain (analyze, costs off) select count(*) from jr where  
jr @@ 'similar_product_ids && ["B000089778"];
```

QUERY PLAN

Aggregate (actual time=0.359..0.359 rows=1 loops=1)

-> Bitmap Heap Scan on jr (actual time=0.084..0.337 rows=185 loops=1)

Recheck Cond: (jr @@ "similar_product_ids" && ["B000089778"]::jsquery)

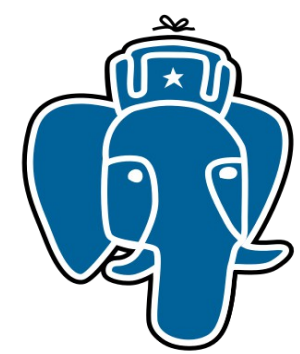
Heap Blocks: exact=107

-> Bitmap Index Scan on jr_path_value_idx (actual time=0.057..0.057 rows=185 loops=1)

Index Cond: (jr @@ "similar_product_ids" && ["B000089778"]::jsquery)

Execution time: 0.394 ms

(7 rows)



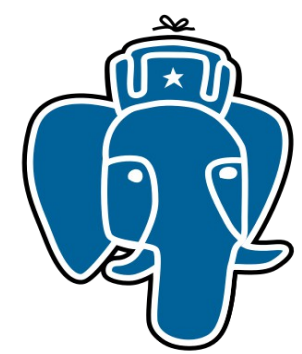
Jsquery (indexes)

- No statistics, no planning :(

```
explain (analyze, costs off) select count(*) from jr where
jr @@ 'similar_product_ids && ["B000089778"]
AND product_sales_rank($ > 10000 AND $ < 20000)';
QUERY PLAN
```

Not selective, better not use index!

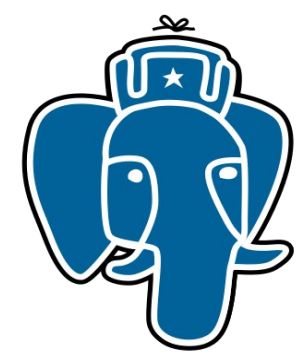
```
Aggregate (actual time=126.149..126.149 rows=1 loops=1)
-> Bitmap Heap Scan on jr (actual time=126.057..126.143 rows=45 loops=1)
    Recheck Cond: (jr @@ ("similar_product_ids" && ["B000089778"]) &
"product_sales_rank"($ > 10000 & $ < 20000))::jsquery)
    Heap Blocks: exact=45
-> Bitmap Index Scan on jr_path_value_idx (actual time=126.029..126.029 rows=45 loops=1)
    Index Cond: (jr @@ ("similar_product_ids" && ["B000089778"]) &
"product_sales_rank"($ > 10000 & $ < 20000))::jsquery)
Execution time: 129.309 ms !!! No statistics
(7 rows)
```



MongoDB 2.6.0

```
db.reviews.find( { $and : [ {similar_product_ids: { $in:["B000089778"]}}, {product_sales_rank:{$gt:10000, $lt:20000}} ] } )
.explain()
{
  "n" : 45,
  .....
  "millis" : 7,
  "indexBounds" : {
    "similar_product_ids" : [
      [
        "B000089778",
        "B000089778"
      ]
    ]
  },
}
```

index size = 400 MB just for similar_product_ids !!!



Jsquery (indexes)

- If we rewrite query and use planner

```
explain (analyze, costs off) select count(*) from jr where  
jr @@ 'similar_product_ids && ["B000089778"]'
```

```
and (jr->>'product_sales_rank')::int>10000 and (jr->>'product_sales_rank')::int<20000;
```

```
Aggregate (actual time=0.479..0.479 rows=1 loops=1)
```

```
-> Bitmap Heap Scan on jr (actual time=0.079..0.472 rows=45 loops=1)
```

```
Recheck Cond: (jr @@ "similar_product_ids" && ["B000089778"]>::jsquery)
```

```
Filter: (((jr ->> 'product_sales_rank'::text))::integer > 10000) AND
```

```
((jr ->> 'product_sales_rank'::text))::integer < 20000))
```

```
Rows Removed by Filter: 140
```

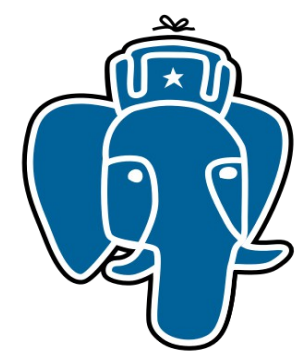
```
Heap Blocks: exact=107
```

```
-> Bitmap Index Scan on jr_path_value_idx (actual time=0.041..0.041 rows=185 loops=1)
```

```
Index Cond: (jr @@ "similar_product_ids" && ["B000089778"]>::jsquery)
```

```
Execution time: 0.506 ms Potentially, query could be faster Mongo !
```

```
(9 rows)
```

Jsquery (optimiser) — **NEW !**

- Jsquery now has built-in optimiser for simple queries.

```
explain (analyze, costs off) select count(*) from jr where  
jr @@ 'similar_product_ids && ["B000089778"]  
AND product_sales_rank( $ > 10000 AND $ < 20000)'
```

```
Aggregate (actual time=0.422..0.422 rows=1 loops=1)
```

```
-> Bitmap Heap Scan on jr (actual time=0.099..0.416 rows=45 loops=1)
```

```
Recheck Cond: (jr @@ '("similar_product_ids" && ["B000089778"]) AND  
"product_sales_rank"($ > 10000 AND $ < 20000))::jsquery)
```

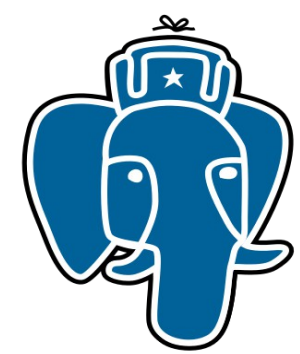
```
Rows Removed by Index Recheck: 140
```

```
Heap Blocks: exact=107
```

```
-> Bitmap Index Scan on jr_path_value_idx (actual time=0.060..0.060 rows=185 loops=1)
```

```
Index Cond: (jr @@ '("similar_product_ids" && ["B000089778"]) AND  
"product_sales_rank"($ > 10000 AND $ < 20000))::jsquery)
```

```
Execution time: 0.480 ms vs 7 ms MongoDB !
```



Jsquery (optimiser) — NEW !

- Jsquery now has built-in optimiser for simple queries.
Analyze query tree and push non-selective parts to recheck (like filter)

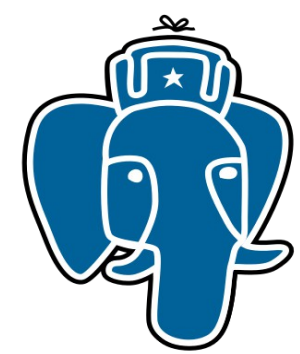
Selectivity classes:

- 1) Equality ($x = c$)
- 2) Range ($c1 < x < c2$)
- 3) Inequality ($c > c1$)
- 4) Is (x is type)
- 5) Any ($x = *$)

```
SELECT gin_debug_query_path_value('similar_product_ids && ["B000089778"]  
AND product_sales_rank( $ > 10000 AND $ < 20000)');
```

```
gin_debug_query_path_value
```

```
-----  
similar_product_ids.# = "B000089778" , entry 0 +
```



Jsquery (optimiser) — NEW !

- Jsquery optimiser pushes non-selective operators to recheck

```
explain (analyze, costs off) select count(*) from jr where  
jr @@ 'similar_product_ids && ["B000089778"]  
AND product_sales_rank( $ > 10000 AND $ < 20000)'
```

```
Aggregate (actual time=0.422..0.422 rows=1 loops=1)
```

```
-> Bitmap Heap Scan on jr (actual time=0.099..0.416 rows=45 loops=1)
```

```
Recheck Cond: (jr @@ ("similar_product_ids" && ["B000089778"]) AND
```

```
"product_sales_rank"($ > 10000 AND $ < 20000))::jsquery)
```

```
Rows Removed by Index Recheck: 140
```

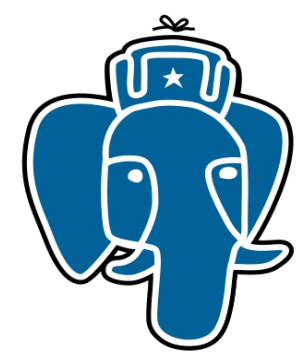
```
Heap Blocks: exact=107
```

```
-> Bitmap Index Scan on jr_path_value_idx (actual time=0.060..0.060 rows=185 loops=1)
```

```
Index Cond: (jr @@ ("similar_product_ids" && ["B000089778"]) AND
```

```
"product_sales_rank"($ > 10000 AND $ < 20000))::jsquery)
```

```
Execution time: 0.480 ms
```



Jsquery (HINTING) — NEW !

- Jsquery now has HINTING (if you don't like optimiser)!

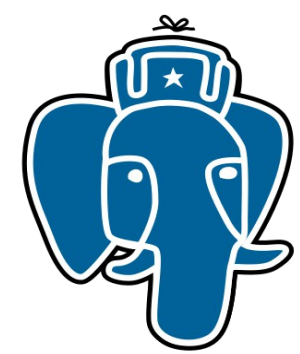
```
explain (analyze, costs off) select count(*) from jr where jr @@ 'product_sales_rank > 10000'
```

```
-----  
Aggregate (actual time=2507.410..2507.410 rows=1 loops=1)  
-> Bitmap Heap Scan on jr (actual time=1118.814..2352.286 rows=2373140 loops=1)  
    Recheck Cond: (jr @@ "product_sales_rank" > 10000)::jsquery)  
    Heap Blocks: exact=201209  
-> Bitmap Index Scan on jr_path_value_idx (actual time=1052.483..1052.48  
rows=2373140 loops=1)  
    Index Cond: (jr @@ "product_sales_rank" > 10000)::jsquery)  
Execution time: 2524.951 ms
```

- Better not to use index — HINT /* --noindex */

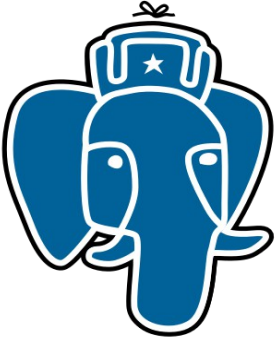
```
explain (analyze, costs off) select count(*) from jr where jr @@ 'product_sales_rank /*-- noindex */ > 10000';
```

```
-----  
Aggregate (actual time=1376.262..1376.262 rows=1 loops=1)  
-> Seq Scan on jr (actual time=0.013..1222.123 rows=2373140 loops=1)  
    Filter: (jr @@ "product_sales_rank" /*-- noindex */ > 10000)::jsquery)  
    Rows Removed by Filter: 650022  
Execution time: 1376.284 ms
```

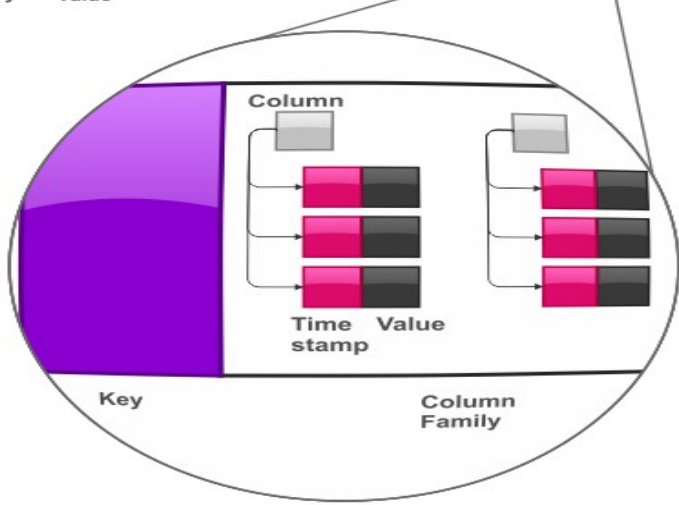
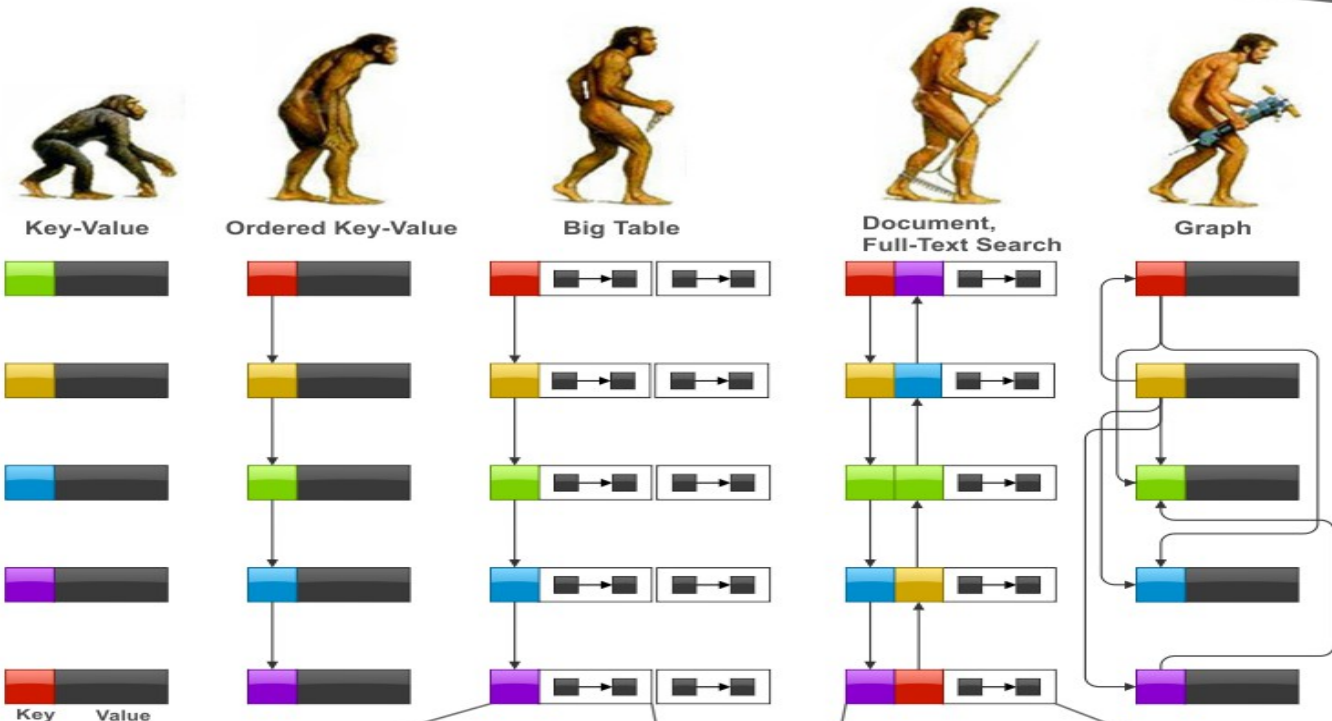
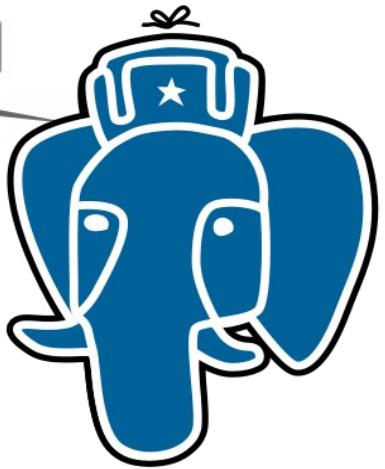


Contrib/jsquery

- Jsquery index support is quite efficient (0.5 ms vs Mongo 7 ms !)
- Future direction
 - Make jsquery planner friendly
 - Need statistics for jsonb
- Availability
 - Jsquery + opclasses are available as extensions
 - Grab it from <https://github.com/akorotkov/jsquery> (branch master) , we need your feedback !
 - We will release it after PostgreSQL 9.4 release
 - Need real sample data and queries !



Stop following me, you fucking freaks!

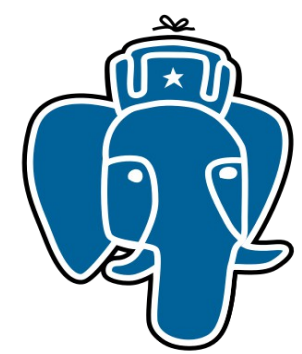


```

employee" :
{
  "name" : "Mohana Pillai"
  "position" : "Delivery Manager"
  "projects" : [
    {
      "name" : "Easy Signu
    }
  ], Semi-Structured Data
}
Plain Text
is a confidential word or number
combination used as a code to
identity when accessing
en 8 and 15 characters
number and may ne
spaces

```

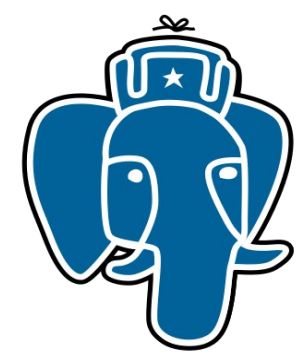
- PostgreSQL 9.4+
- Open-source
 - Relational database
 - Strong support of json



Better indexing ...

- GIN is a proven and effective index access method
- Need indexing for jsonb with operations on paths (no hash!) and values
 - B-tree in entry tree is not good - length limit, no prefix compression

List of relations							
Schema	Name	Type	Owner	Table	Size	Description	
public	jb	table	postgres		1374 MB		
public	jb_uniq_paths	table	postgres		912 MB		
public	jb_uniq_paths_btree_idx	index	postgres	jb_uniq_paths	885 MB	text_pattern_ops	
public	jb_uniq_paths_spgist_idx	index	postgres	jb_uniq_paths	598 MB	now much less !	

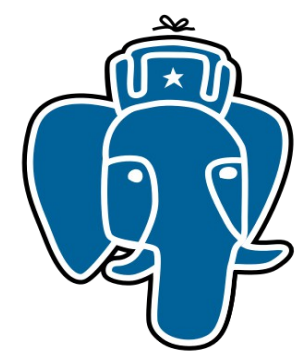


Better indexing ...

- Provide interface to change hardcoded B-tree in Entry tree
 - Use spgist opclass for storing paths and values as is (strings hashed in values)
- We may go further - provide interface to change hardcoded B-tree in posting tree
 - GIS aware full text search !
- New index access method

CREATE INDEX ... USING VODKA





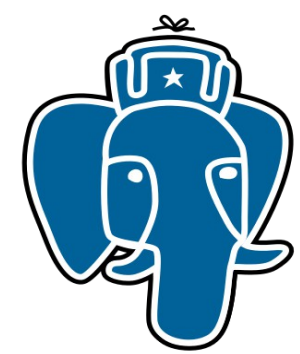
GIN History

- Introduced at PostgreSQL Anniversary Meeting in Toronto, Jul 7-8, 2006 by Oleg Bartunov and Teodor Sigaev



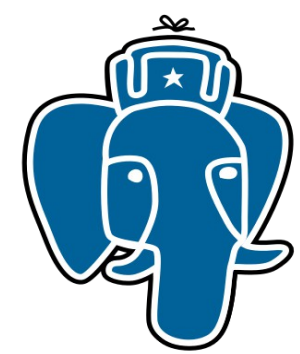
Generalized Inverted Index

- An inverted index is an index structure storing a set of (key, posting list) pairs, where 'posting list' is a set of documents in which the key occurs.
- Generalized means that the index does not know which operation it accelerates. It works with custom strategies, defined for specific data types. GIN is similar to GiST and differs from B-Tree indices, which have predefined, comparison-based operations.



GIN History

- Introduced at PostgreSQL Anniversary Meeting in Toronto, Jul 7-8, 2006 by Oleg Bartunov and Teodor Sigaev
- Supported by JFG Networks (France)
- «Gin stands for Generalized Inverted iNdex and should be considered as a genie, not a drink.»
- Alexander Korotkov, Heikki Linnakangas have joined GIN++ development in 2013



GIN History

- From GIN Readme, posted in -hackers, 2006-04-26

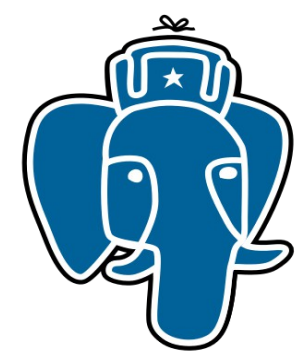
TODO

Nearest future:

- * Opclasses for all types (no programming, just many catalog changes).

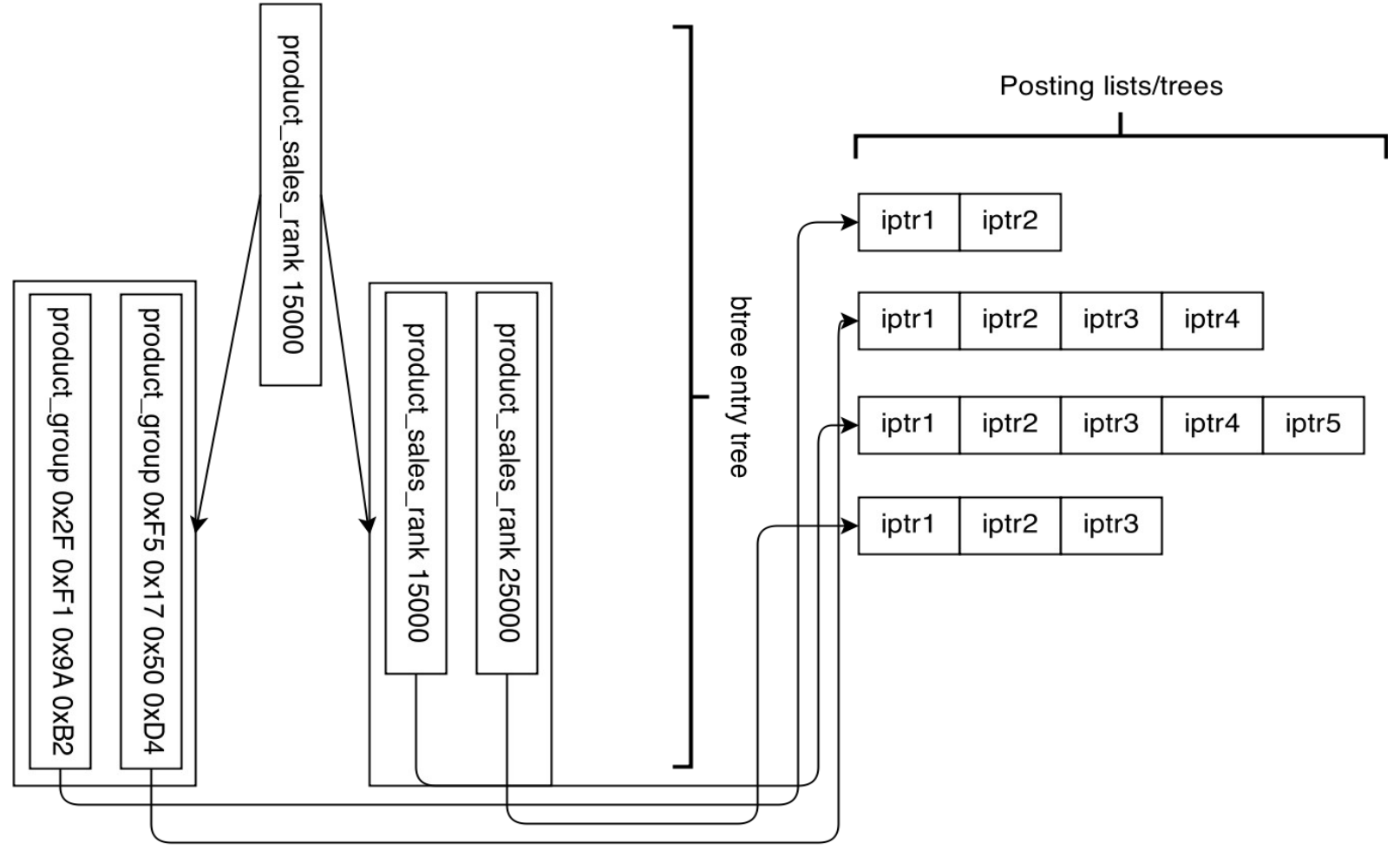
Distant future:

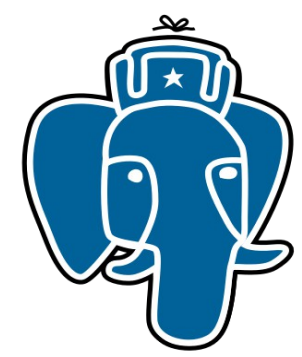
- * Replace B-tree of entries to something like GiST (**VODKA ! 2014**)
- * Add multicolumn support
- * Optimize insert operations (background index insertion)



GIN index structure for jsonb

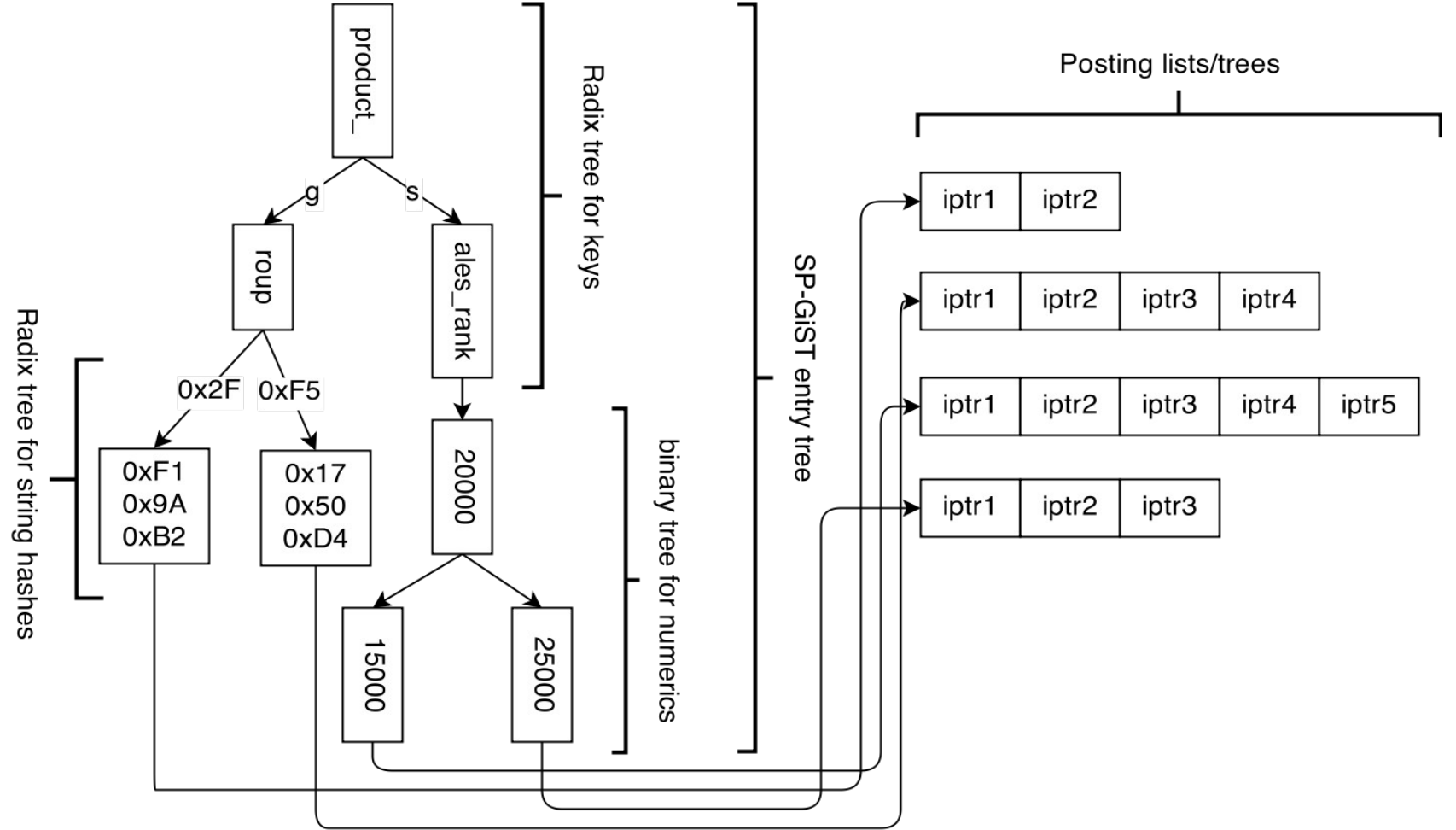
```
{  
  "product_group": "Book",  
  "product_sales_rank": 15000  
},  
{  
  "product_group": "Music",  
  "product_sales_rank": 25000  
}
```

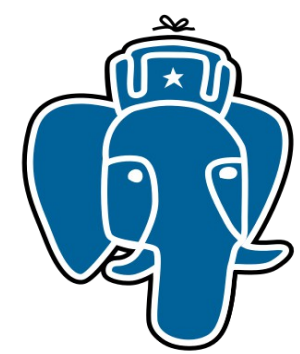




Vodka index structure for jsonb

```
{  
  "product_group": "Book",  
  "product_sales_rank": 15000  
},  
{  
  "product_group": "Music",  
  "product_sales_rank": 25000  
}
```





CREATE INDEX ... USING VODKA

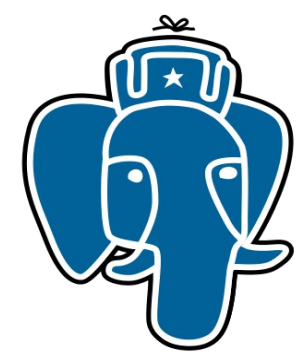
- Delicious bookmarks, mostly text data

```
set maintenance_work_mem = '1GB';
```

List of relations

Schema	Name	Type	Owner	Table	Size	Description
public	jb	table	postgres		1374 MB	1252973 rows
public	jb_value_path_idx	index	postgres	jb	306 MB	98769.096
public	jb_gin_idx	index	postgres	jb	544 MB	129860.859
public	jb_path_value_idx	index	postgres	jb	306 MB	100560.313
public	jb_path_idx	index	postgres	jb	251 MB	68880.320
public	jb_vodka_idx	index	postgres	jb	409 MB	185362.865
public	jb_vodka_idx5	index	postgres	jb	325 MB	174627.234 new spgist

(6 rows)



CREATE INDEX ... USING VODKA

```
select count(*) from jb where jb @@ 'tags.#.term' = "NYC";
```

Aggregate (actual time=0.423..0.423 rows=1 loops=1)

-> Bitmap Heap Scan on jb (actual time=0.146..0.404 rows=285 loops=1)

Recheck Cond: (jb @@ "tags".#"term" = "NYC"::jsquery)

Heap Blocks: exact=285

-> Bitmap Index Scan on jb_vodka_idx (actual time=0.108..0.108 rows=285 loops=1)

Index Cond: (jb @@ "tags".#"term" = "NYC"::jsquery)

Execution time: 0.456 ms (0.634 ms, GIN jsonb_value_path_ops)

```
select count(*) from jb where jb @@ '*.term' = "NYC";
```

Aggregate (actual time=0.495..0.495 rows=1 loops=1)

-> Bitmap Heap Scan on jb (actual time=0.245..0.474 rows=285 loops=1)

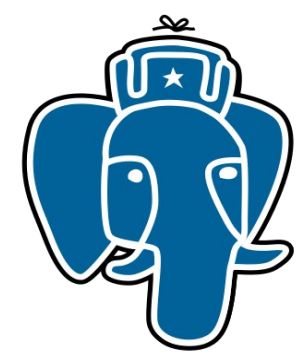
Recheck Cond: (jb @@ '*."term' = "NYC"::jsquery)

Heap Blocks: exact=285

-> Bitmap Index Scan on jb_vodka_idx (actual time=0.214..0.214 rows=285 loops=1)

Index Cond: (jb @@ '*."term' = "NYC"::jsquery)

Execution time: 0.526 ms (0.716 ms, GIN jsonb_path_value_ops)



CREATE INDEX ... USING VODKA

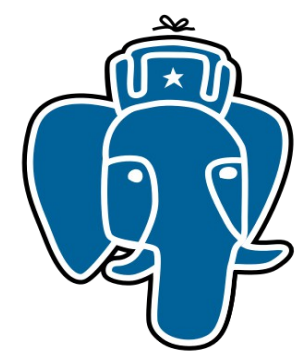
- CITUS data, text and numeric

```
set maintenance_work_mem = '1GB';
```

List of relations

Schema	Name	Type	Owner	Table	Size	Description
public	jr	table	postgres		1573 MB	3023162 rows
public	jr_value_path_idx	index	postgres	jr	196 MB	79180.120
public	jr_gin_idx	index	postgres	jr	235 MB	111814.929
public	jr_path_value_idx	index	postgres	jr	196 MB	73369.713
public	jr_path_idx	index	postgres	jr	180 MB	48981.307
public	jr_vodka_idx3	index	postgres	jr	240 MB	155714.777
public	jr_vodka_idx4	index	postgres	jr	211 MB	169440.130 new spgist

(6 rows)



CREATE INDEX ... USING VODKA

```
explain (analyze, costs off) select count(*) from jr where jr @@ 'similar_product_ids && ["B000089778"]';  
QUERY PLAN
```

Aggregate (actual time=0.200..0.200 rows=1 loops=1)

-> Bitmap Heap Scan on jr (actual time=0.090..0.183 rows=185 loops=1)

Recheck Cond: (jr @@ "similar_product_ids" && ["B000089778"]::jsquery)

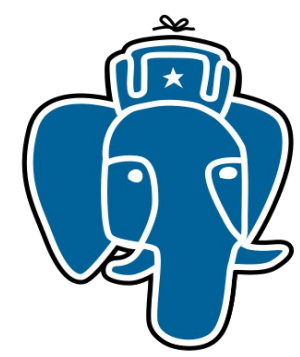
Heap Blocks: exact=107

-> Bitmap Index Scan on jr_vodka_idx (actual time=0.077..0.077 rows=185 loops=1)

Index Cond: (jr @@ "similar_product_ids" && ["B000089778"]::jsquery)

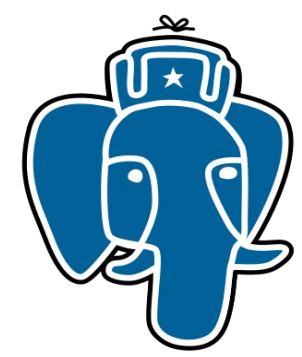
Execution time: 0.237 ms (0.394 ms, GIN jsonb_path_value_idx)

(7 rows)



There are can be different flavors of Vodka

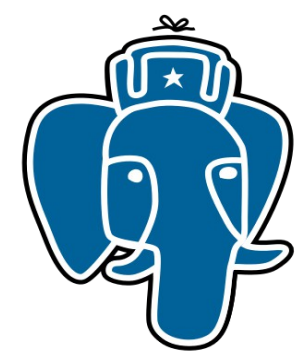




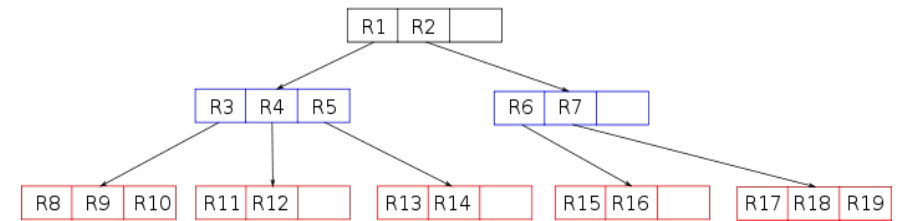
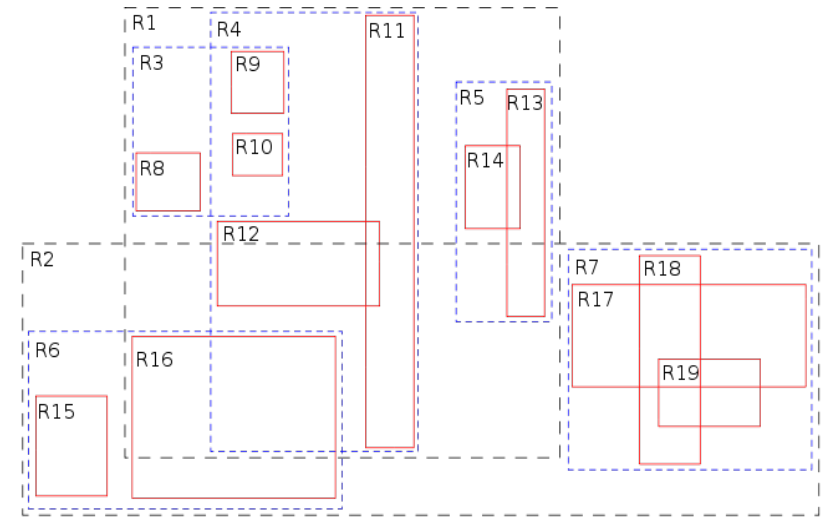
Spaghetti indexing ...



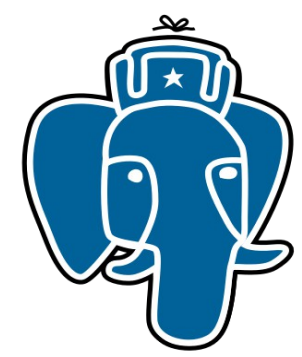
Find twirled spaghetti



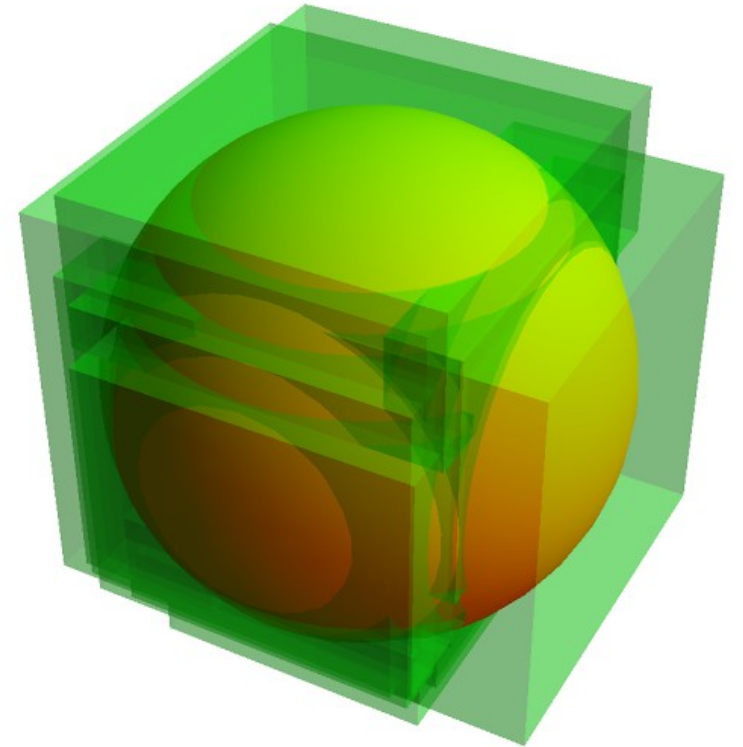
Spaghetti indexing ...



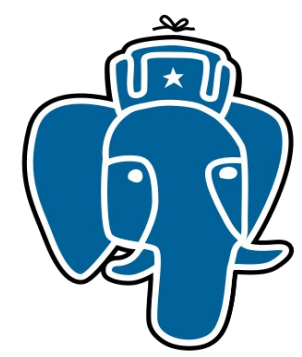
R-tree fails here – bounding box of each separate spaghetti is the same



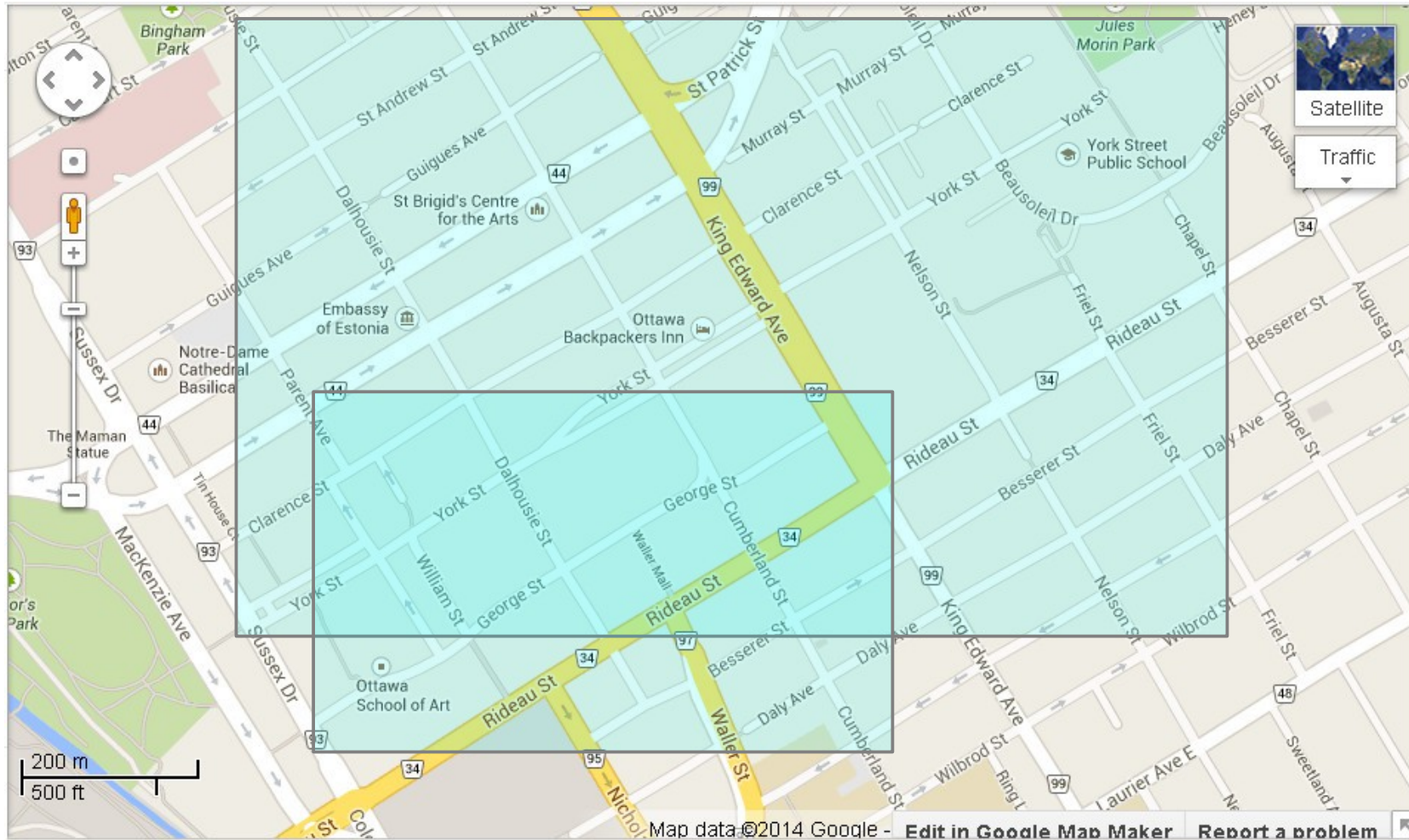
Spaghetti indexing ...

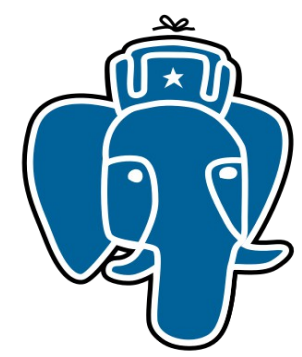


R-tree fails here — bounding box of each separate spaghetti is the same



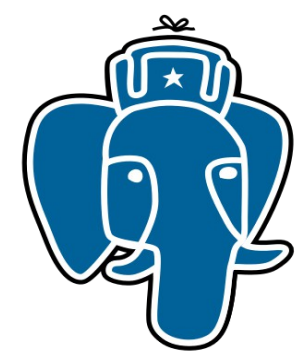
Ottawa downtown: York and George streets



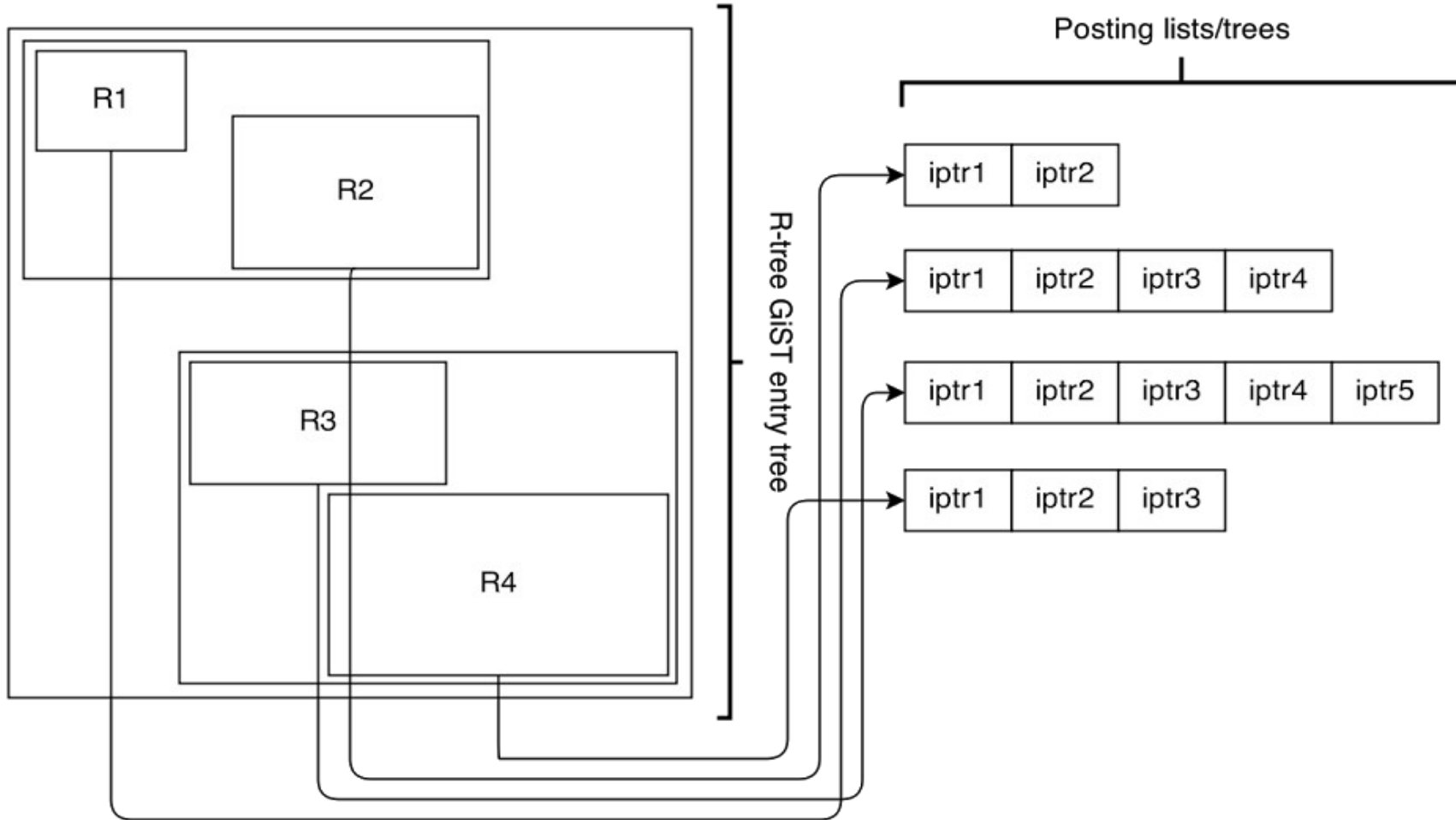


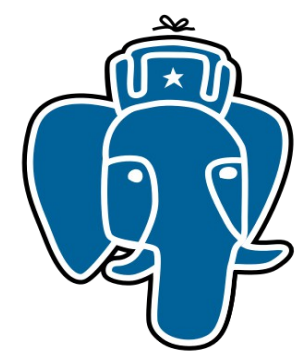
Idea: Use multiple boxes






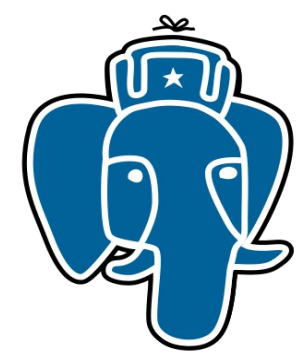
Rtree Vodka





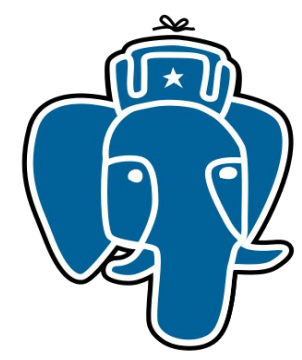
Summary

- contrib/jsquery for 9.4
 - Jsquery - Jsonb Query Language
 - Two GIN opclasses with jsquery support
 - Grab it from <https://github.com/akorotkov/jsquery> (branch master)
- Prototype of VODKA access method
- Plans for improving indexing infrastructure
- This work was supported by  heroku

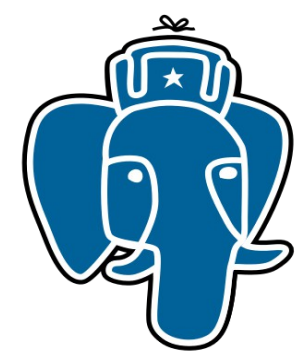


Another view on VODKA

- VODKA CONNECTING INDEXES
 - composite index, which combines different access methods
 - Nested search trees



ご清聴ありがとうございました



VODKA Optimized Dendriform Keys Array